

## PSEUDOCODE FOR THE \$Q RECOGNIZER

We provide complete pseudocode for \$Q. In the following, POINT is a structure that exposes  $x$ ,  $y$ , and  $strokeId$  properties;  $strokeId$  is the stroke index a point belongs to and is filled by counting touch down/touch up events;  $x$  and  $y$  are integers in the set  $\{0, 1, \dots, m-1\}$ , where  $m$  is the size of the look-up table. POINTS is a list of points and TEMPLATES a list of POINTS with associated gesture class data. The pseudocode assumes that templates have already been preprocessed (when loading the template set, for instance).

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### \$Q-RECOGNIZER (POINTS $points$ , TEMPLATES $templates$ )

```

1:  $n \leftarrow 32, m \leftarrow 64$  // defaults for cloud size ( $n$ ) and size of the look-up table ( $m$ )
2: NORMALIZE( $points, n, m$ ) // templates have already been normalized
3:  $score \leftarrow \infty$ 
4: for each  $template$  in  $templates$  do
5:    $d \leftarrow$  CLOUD-MATCH( $points, template, n, score$ )
6:   if  $d < score$  then
7:      $score \leftarrow d$ 
8:    $result \leftarrow template$ 
9: return ( $result, score$ ) // the closest  $template$  from the set and the smallest  $score$ 

```

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### CLOUD-MATCH (POINTS $points$ , POINTS $template$ , int $n$ , int $min$ )

```

1:  $step \leftarrow \lfloor n^{0.5} \rfloor$ 
2: // compute lower bounds for both matching directions between  $points$  and  $template$ 
3:  $LB_1 \leftarrow$  COMPUTE-LOWER-BOUND( $points, template, step, template.LUT$ )
4:  $LB_2 \leftarrow$  COMPUTE-LOWER-BOUND( $template, points, step, points.LUT$ )
5: for  $i \leftarrow 0$  to  $n-1$  step  $step$  do
6:   if  $LB_{1\lfloor i/step \rfloor} < min$  then
7:      $min \leftarrow$  MIN( $min, CLOUD-DISTANCE(points, template, n, i, min)$ )
8:   if  $LB_{2\lfloor i/step \rfloor} < min$  then
9:      $min \leftarrow$  MIN( $min, CLOUD-DISTANCE(template, points, n, i, min)$ )
10: return  $min$ 

```

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### CLOUD-DISTANCE (POINTS $points$ , POINTS $template$ , int $n$ , int $start$ , float $minSoFar$ )

```

1:  $unmatched \leftarrow \{0, 1, 2, \dots, n-1\}$  // indices of unmatched points from  $template$ 
2:  $i \leftarrow start$  // start the matching from this index in the  $points$  cloud
3:  $weight \leftarrow n$  // weights decrease from  $n$  to 1
4:  $sum \leftarrow 0$  // computes the cloud distance between  $points$  and  $template$ 
5: do
6:    $min \leftarrow \infty$ 
7:   for each  $j$  in  $unmatched$  do
8:      $d \leftarrow$  SQR-EUCLIDEAN-DISTANCE( $points_{s_{[j]}}$ ,  $template_{[j]}$ )
9:     if  $d < min$  then
10:        $min \leftarrow d$ 
11:        $index \leftarrow j$ 
12: REMOVE( $unmatched, index$ ) // implementable in  $O(1)$ 
13:  $sum \leftarrow sum + weight \cdot min$ 
14: if  $sum \geq minSoFar$  then
15:   return  $sum$  // early abandoning of computations
16:  $weight \leftarrow weight - 1$  // weights decrease from  $n$  to 1
17:  $i \leftarrow (i+1) \text{ MOD } n$  // advance to the next point in  $points$ 
18: until  $i == start$ 
19: return  $sum$ 

```

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### COMPUTE-LOWER-BOUND (POINTS $points$ , POINTS $template$ , int $step$ , int[,] $LUT$ )

```

1:  $LB \leftarrow$  new float[ $n/step + 1$ ] // multiple lower bounds, one for each starting point
2:  $SAT \leftarrow$  new float[ $n$ ] // summed area table for fast computations (see text)
3: // first, compute the lower bound for starting point index 0
4:  $LB_{[0]} \leftarrow 0$ 
5: for  $i \leftarrow 0$  to  $n-1$  do
6:    $index \leftarrow LUT[points_{s_{[i]}.x}, points_{s_{[i]}.y}]$ 
7:    $d \leftarrow$  SQR-EUCLIDEAN-DISTANCE( $points_{s_{[i]}}$ ,  $template_{[index]}$ )
8:    $SAT_{[i]} \leftarrow (i == 0) ? d : SAT_{[i-1]} + d$ 
9:    $LB_{[0]} \leftarrow LB_{[0]} + (n-i) \cdot d$ 
10: // compute the lower bound for the other starting points (see formula in the text)
11: for  $i \leftarrow step$  to  $n-1$  step  $step$  do
12:    $LB_{\lfloor i/step \rfloor} \leftarrow LB_{[0]} + i \cdot SAT_{[n-1]} - n \cdot SAT_{[i-1]}$ 
13: return  $LB$ 

```

The following pseudocode implements gesture preprocessing: resampling, translation to origin, rescaling into the  $m \times m$  grid, and computation of the look-up table. Except for the new COMPUTE-LUT function and changes in the SCALE function, this pseudocode is practically the same as for \$P [50] (p. 280).

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### NORMALIZE (POINTS $points$ , int $n$ , int $m$ )

```

1:  $points \leftarrow$  RESAMPLE( $points, n$ )
2: TRANSLATE-TO-ORIGIN( $points, n$ )
3: SCALE( $points, m$ )
4: LUT  $\leftarrow$  COMPUTE-LUT( $m, points, n$ )

```

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### RESAMPLE (POINTS $points$ , int $n$ )

```

1:  $I \leftarrow$  PATH-LENGTH( $points$ ) / ( $n-1$ )
2:  $D \leftarrow 0$ 
3:  $newPoints \leftarrow \{points_{[0]}\}$ 
4: for  $i \leftarrow 1$  to  $n-1$  do
5:   if  $points_{[i]}.strokeId == points_{[i-1]}.strokeId$  then
6:      $d \leftarrow$  EUCLIDEAN-DISTANCE( $points_{[i-1]}$ ,  $points_{[i]}$ )
7:     if  $(D+d) \geq I$  then
8:        $q.x \leftarrow points_{[i-1].x} + (I-D)/d \cdot (points_{[i].x} - points_{[i-1].x})$ 
9:        $q.y \leftarrow points_{[i-1].y} + (I-D)/d \cdot (points_{[i].y} - points_{[i-1].y})$ 
10:      APPEND( $newPoints, q$ )
11:      INSERT( $points, i, q$ ) //  $q$  will be the next  $points_{[i]}$ 
12:       $D \leftarrow 0$ 
13:     else  $D \leftarrow D+d$ 
14: return  $newPoints$ 

```

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### TRANSLATE-TO-ORIGIN (POINTS $points$ , int $n$ )

```

1:  $c \leftarrow (0,0)$  // will compute the centroid of the  $points$  cloud
2: for each  $p$  in  $points$  do
3:    $c \leftarrow (c.x + p.x, c.y + p.y)$ 
4:  $c \leftarrow (c.x/n, c.y/n)$ 
5: for each  $p$  in  $points$  do
6:    $p \leftarrow (p.x - c.x, p.y - c.y)$ 

```

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### SCALE (POINTS $points$ , int $m$ )

```

1:  $x_{min} \leftarrow \infty, x_{max} \leftarrow -\infty, y_{min} \leftarrow \infty, y_{max} \leftarrow -\infty$ 
2: for each  $p$  in  $points$  do
3:    $x_{min} \leftarrow$  MIN( $x_{min}, p.x$ )
4:    $y_{min} \leftarrow$  MIN( $y_{min}, p.y$ )
5:    $x_{max} \leftarrow$  MAX( $x_{max}, p.x$ )
6:    $y_{max} \leftarrow$  MAX( $y_{max}, p.y$ )
7:  $s \leftarrow$  MAX( $x_{max} - x_{min}, y_{max} - y_{min}$ ) / ( $m-1$ ) // scale factor
8: for each  $p$  in  $points$  do
9:    $p \leftarrow ((p.x - x_{min})/s, (p.y - y_{min})/s)$  //  $p.x$  and  $p.y$  are now integers in  $0..m-1$ 

```

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### COMPUTE-LUT (POINTS $points$ , int $n$ , int $m$ )

```

1: LUT  $\leftarrow$  new int[ $m, m$ ]
2: for  $x \leftarrow 0$  to  $m-1$  do
3:   for  $y \leftarrow 0$  to  $m-1$  do
4:      $min \leftarrow \infty$ 
5:     for  $i \leftarrow 0$  to  $n-1$  do
6:        $d \leftarrow$  SQR-EUCLIDEAN-DISTANCE( $points_{[i]}$ , new POINT( $x, y$ ))
7:       if  $d < min$  then
8:          $min \leftarrow d$ 
9:          $index \leftarrow i$ 
10:      LUT[ $x, y$ ]  $\leftarrow index$ 
11: return LUT

```

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### PATH-LENGTH (POINTS $points$ )

```

1:  $d \leftarrow 0$  // will compute the path length
2: for  $i \leftarrow 1$  to  $n-1$  do
3:   if  $points_{[i]}.strokeId == points_{[i-1]}.strokeId$  then
4:      $d \leftarrow d +$  EUCLIDEAN-DISTANCE( $points_{[i-1]}$ ,  $points_{[i]}$ )
5: return  $d$ 

```

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### SQR-EUCLIDEAN-DISTANCE (POINT $a$ , POINT $b$ )

```

1: return  $(a.x - b.x)^2 + (a.y - b.y)^2$  // much faster to compute without the  $\sqrt{\quad}$ ()

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### EUCLIDEAN-DISTANCE (POINT $a$ , POINT $b$ )

```

1: return  $\sqrt{\text{SQR-EUCLIDEAN-DISTANCE}(a, b)}$ 

```