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Hilbert schemes

Want to study moduli spaces, i.e.
some space $\mathcal{M}$ so that

$$\text{Hom}(B, \mathcal{M}) = \left\{ \begin{array}{l} \text{flat families of} \\ \text{some kind of objects}/B \end{array} \right\}.$$

It is not clear how to go about
constructing such $\mathcal{M}$ (fine or
coarse)

General strategy: Add some extra data
to your objects, and

(1) find a fine moduli space
for (objects + extra data).

(2) "quotient" the answer from
(1) by a suitable equivalence
relation killing the extra data

The Hilbert scheme is a solution to
(1). It is a fine moduli space
and a projective scheme

We will study the moduli problem describing the deformations of a projective variety X .

Extra data: Instead of considering X as an abstract variety look at X , along with an embedding

$$X \hookrightarrow \mathbb{P}^n$$

How does the presence of the extra data modify our moduli problem?

We want to look at flat families of objects in $\mathbb{P}^n$ that "look like X "

We need to be more precise about what "looks like X " means.

To do this we use the Hilbert polynomial.

Two definitions of Hilbert polynomials

(I) (Hilbert) Let $A = k[z_0, \dots, z_n]$

be a polynomial ring, and let

$$I_X \triangleleft A$$

be the homogeneous ideal defining X .

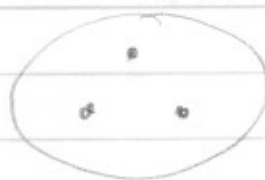
The Hilbert function of X is the function

$$HF_X(d) := \dim_k \underbrace{(A/I_X)_d}_{\uparrow \text{ in degree } d}$$

Thm: For large d , $HF_X(d)$ is given by a polynomial in d .

We will write $P_X(d)$ for that polynomial.

Examples: (1) Suppose $X \subset \mathbb{P}^2$ is the collection of 3 general points.



In this case

d	0	1	2	3	4
HF_X	1	3	3	3	3
P_X	3	3	3	3	3

$$p_X(d) = 3$$

(2) Suppose $X \subset \mathbb{P}^2$ consists of 3 collinear points. In this case

d	0	1	2	3	4
HF_X	1	2	3	3	3
P_X	3	3	3	3	3

$$p_X(d) = 3$$

(3) $X =$ intersection of two degree 4 hypersurfaces in $\mathbb{P}^3$

In the transversal case X is a smooth curve

$$\deg(X) = 16$$

$$g(X) = 33$$

d	0	1	2	3	4	5	6	7
HF_X	1	4	10	20	33	48	64	80
P_X	-32	-16	0	16	32	48	64	80

$\xrightarrow{\hspace{10em}}$
 HF_X becomes a poly

$$P_X(t) = 16t - 32$$

Why is the Hilbert function a polynomial for large d ?

Look at example (3). Here we can resolve A/I

$$0 \rightarrow A(-8) \rightarrow \bigoplus_{i=1}^4 A(-4) \rightarrow A \rightarrow (A/I) \rightarrow 0$$

Here $A(-k)$ = shifted version of A

$$[A(-k)]_d := A_{d-k}$$

Now

$$\dim (A/I)_d = \dim A_d - 2 \dim A(-4)_d + \dim A(-8)_d$$

$$= \binom{d+3}{3} - 2 \binom{d-4+3}{3} + \binom{d-8+3}{3}$$

here $\binom{m}{n} = \begin{cases} 0 & \text{if } m < n \\ \text{usual } \binom{m}{n} & \text{if } m \geq n \end{cases}$

Thus HF_X is a poly for large d .

Note: This is the general proof that HF_X is a poly for large d . To carry out the proof in general we need Hilbert's Syzygy theorem - resolutions stop after finitely many steps.

(II) (Serre) let $\mathcal{O}_X(d) = \mathcal{O}_X \otimes_{\mathcal{O}_{\mathbb{P}^n}} \mathcal{O}_{\mathbb{P}^n}(d)$

Define

$$\chi(\mathcal{O}_X(d)) = \sum_{i \geq 0} (-1)^i \dim H^i(X, \mathcal{O}_X(d))$$

Thm This is a polynomial function in $d = P_X(d)$

Sketch of proof: To check that

$\chi(\mathcal{O}_X(d))$ is a polynomial take
 Y - general hyperplane section of
 $X \Rightarrow$ get an exact sequence

$$0 \rightarrow \mathcal{O}_X(-1) \rightarrow \mathcal{O}_X \rightarrow \mathcal{O}_Y \rightarrow 0$$

Twist by $\mathcal{O}_X(d) \Rightarrow$

$$0 \rightarrow \mathcal{O}_X(d-1) \rightarrow \mathcal{O}_X(d) \rightarrow \mathcal{O}_Y(d) \rightarrow 0$$

$\Rightarrow$

$$\chi(\mathcal{O}_Y(d)) = \chi(\mathcal{O}_X(d)) - \chi(\mathcal{O}_X(d-1))$$

polynomial
by induction

the first difference
of $\chi(\mathcal{O}_X(d))$

But if the first difference of
a function is polynomial, then
the function must be polynomial

The fact that $\chi(\mathcal{O}_X(d))$ and
 $P_X(d)$ agree is a consequence
of Serre vanishing.

Consider the short exact sequence

$$0 \rightarrow I_X(d) \rightarrow \mathcal{O}_{\mathbb{P}^n}(d) \rightarrow \mathcal{O}_X(d) \rightarrow 0$$

For large d we have

$$\chi(I_X(d)) = \dim(I_X)_d$$

$$\chi(\mathcal{O}_{\mathbb{P}^n}(d)) = \dim(A)_d$$

$$\chi(\mathcal{O}_X(d)) = \dim(A/I_X)_d$$

by Serre vanishing

□

Going back to the flat families of things that "look like X " in $\mathbb{P}^n$: What does the Hilbert polynomial tell us?

If $\dim X = m$, $\deg X = d$, then

$$P_X(t) = \frac{d}{m!} t^m + \text{lower terms} + \chi(\mathcal{O}_X)$$

Example: If X is a curve of genus g and degree d in

in same $\mathbb{P}^n$, then

$$p_X(t) = dt + (1-g)$$

Similarly, if X is a smooth surface in $\mathbb{P}^n$ and $L = \mathcal{O}_{\mathbb{P}^n}(1)|_X$

then

$$p_X(t) = \frac{c_1(L)^2}{2} t^2 - \frac{c_1(K) c_1(L)}{2} t + \chi(\mathcal{O}_X)$$

If B - connected scheme and

$$X \subset B \times \mathbb{P}^n$$

a family of subschemes in $\mathbb{P}^n$ parameterized by B , i.e.

$$X \subset B \times \mathbb{P}^n$$

$$\begin{array}{c} \searrow \swarrow p_B^* \\ B \end{array}$$

then, for each $b \in B$ we can look at $X_b \subset \mathbb{P}^n$ and its Hilbert polynomial $p_{X_b}(t)$

Thm (Serre) If X is flat over B ,
 then $P_{X_b}(t)$ is independent
 of $b \in B$.

Conversely if $P_{X_b}(t)$ is
 independent of b and B
 is reduced then the family
 is flat.

Now this suggests the following
 setup. Given a polynomial $P(t)$,
 look for a fine moduli scheme
 $\text{Hilb}_{P(t)}$ which parameterizes
 flat families of schemes
 in $\mathbb{P}^n$
 with Hilbert polynomial $P(t)$.

Example: If $\Lambda \subset \mathbb{P}^n$ is a
 k -plane then the Hilbert
 polynomial of Λ is

$$P_{\Lambda}(t) = \binom{t+k}{k}$$

The Hilbert scheme in this case
 is just the Grassmannian which

we know how to construct.

What can we do in general?

Let us try and parameterise all ideals I_X . The essential difficulty is that I_X is infinite dimensional over k so we should make some reductions first.

Remarks: • If we know $I_{X,d}$ (i.e. the ideal in degree d), then we know $I_{X,d'}$ for all $d' \leq d$.

Indeed if $f \in I_{X,d+1}$, then $z_0 f, z_1 f, \dots, z_n f \in I_{X,d}$ and the common zeros of these are the same zeros of f .

• If I_X is generated in degree $\leq d$, then knowing $I_{X,d}$ lets us compute $I_{X,d'}$ for all d' .

Idea: look at the finite dimensional vector space A_d and the subspace

$$I_{X,d} \subset A_d$$

The dimension $\dim I_{X,d} = ?$

For d large enough we have

$$\dim (A/I_X)_d = P_X(d)$$

$\Rightarrow$ for large d we have

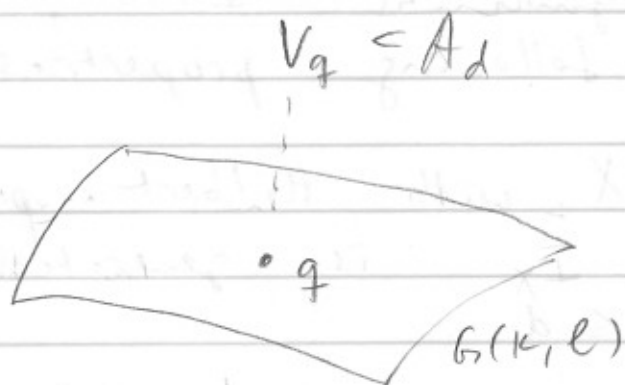
$$\dim I_{X,d} = \binom{n+d}{n} - P_X(d).$$

Now to get the construction going we need to show that we can choose d (large enough) which works for all X 's with our fixed Hilbert polynomial.

Suppose we can find such a d . What more do we need to do to construct $\text{Hilb}_P(t)$?

Fix d . Set $k = \dim I_{X,d}$
 $l = \dim A_d$

Consider the Grassmannian $G(k, l)$
 We need to restrict ourselves to
 points of $G(k, l)$ which generate
 ideals with our Hilbert polynomial



Require that

$$\dim \text{image}(V_q \cdot A_r) \subseteq A_{d+r}$$

is the dimension we expect i.e.

$$= \binom{d+r+n}{n} - R(d+r)$$

There are two issues with this
 setup:

- the property that a linear map has a fixed rank is not a closed condition.

• When do we stop adding equations? (i.e. up to what degree do we need to check?)

Theorem (Gutzmann)

Hilbert polynomial with the following

Given a $P(t)$, $\exists d$ properties

• $\forall X$ with Hilbert poly $P(t) \Rightarrow I_X$ is generated in degree $\leq d$

• $\dim I_{X,d'} = \binom{n+d'}{n} - P(d')$

for all $d' \geq d$.

• for any subspace of dimension $\binom{n+d}{n} - P(d)$ in A_d , the

image of multiplication by $z_0, z_1, \dots, z_n$ in degree $d+1$ is at least

$$\binom{d+1+n}{n} - P(d)$$

In particular, the equations for $d \rightarrow d+1$ are sufficient.

In fact the theorem tells us how to compute d .

Given $P(t)$ there exists a unique sequence of integers

$$a_0 \geq a_1 \geq \dots \geq a_r \geq 0$$

so that

$$P(t) = \binom{t+a_0}{a_0} + \binom{t+a_1}{a_1} + \dots + \binom{t+a_r}{a_r}$$

Note : $\binom{t+a_k-k}{a_k} = \frac{1}{a_k!} t^{a_k} + \dots$

$\Rightarrow d = r+1$ in the theorem

Example: $P(t) = 3t+1$

$$a_0 = 1, \quad a_1 = 1, \quad a_2 = 1, \quad a_3 = 0$$

$$\Rightarrow d = 4.$$