M++ User's Manual
(version 1.00)

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1 Preface

M++ is a new cluster computing paradigm for multi-agent applications viewed as the interactions among autonomous objects, each behaving with its own goals and reacting to the local information in a synthetic world. M++ self-migrating threads are C++ compiled objects, capable of constructing logical computational networks and autonomously migrating themselves over a cluster system. They represent a natural mechanism of implementing multi-agent applications: each individual agent can be represented as a distinct member of collection of self-migrating threads, moving through a logical network representing a synthetic world and interacting with one another. Several actual applications M++ assumes are:

- *Ecological simulation:* simulate a transient population of preys and predators such as fish and sharks [Dew84], and find the genetic code of the most efficient ant forager through a genetic algorithm having many ants searching for food in a farm [CJ92].
- *Artificial life:* simulate the evolution and learning process of neural network for solving a given problem [NNKdG98], and observe the growth of artificial plants [PL96].
- *Complex business modeling:* decide the best cargo-routing strategy by finding which local interactions lead to global behaviors in cargo operations [Wak01].

For more details on M++ suitability to those multi-agent applications, please refer to the following three research papers: [SFB99], [FSCK01], and [LMF$^+$02]. Also, you can obtain further acquaintance with multi-agent applications through the readings: [Fer99] and [Wei99].

This manual provides the procedures necessary for users to start programming M++ self-migrating threads, (simply called M++ threads in the following explanation.) To run M++ threads, you need to install the system, start an M++ daemon at each PC or workstation, code your M++ threads in the M++ language, compile them, and inject the executables to one of active daemons. The manual introduces system installation, daemon invocation, the M++ language specification, thread compilation, and trouble shooting. It also includes code examples to prompt users' understanding.

We welcome any comments and suggestions regarding M++, please contact us at:

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2 Porting and Accessing the System

2.1 Availability

The currently available platforms to run M++ include:

Hardware	Operating Systems
PC	Intel Solaris7 and 8, Redhat Linux 7.3
Sparc	SunOS 5.5.1-SunOS 5.8
AP3000	SunOS 5.6

Table 1: Available Platforms

We have installed M++ at the following sites. Note that *** in Table 2 means a gateway's or each cluster node's IP name. This manual does not disclose this information for security reasons. To find the appropriate name, ask your local system owner or manager. If you are a user at any of those sites, you do not have to install the system. Please skip to Section 2.2.3 to set up your shell environment.

Sites	Platform	OS	IPs and System Managers
UW Bothell	PCs	Redhat Linux	medusa.bothell.washington.edu
			dgrimmer@uwb.edu
Univ. of Tsukuba	PCs	Solaris8	howdy00.padc.mmpc.is.tsukuba.ac.jp
			nas@is.tsukuba.ac.jp
UC Irvine	Sparc	SunOS 5.5.1-5.8	***.ics.uci.edu
			support@ics.uci.edu
Ehime University	AP3000	SunOS 5.6	***.cs.ehime-u.ac.jp
			kob@cs.ehime-u.ac.jp

Table 2: Available Sites

To install M++ on the platforms listed in Table 1, please follow the instructions given in Section 2.2. If you are interested in porting to the other platforms, please refer to Section 2.3.

2.2 Installing M++

First of all, please make sure that your cluster is NFS-mounted. The same disk image must be visible to all computing nodes of the cluster. In addition, each computing node must have an independent IP address, (whether it is a private or a public IP). Unless these requirements are satisfied, the installation is impossible.

2.2.1 Downloading

To obtain a gzipped tar file of M++, visit any of the following M++ web pages:

Sites	Web Addresses
UW Bothell	http://faculty.washington.edu/mfukuda/m++
Univ. of Tsukuba	http://www.padc.mmpc.is.tsukuba.ac.jp/member/naoya/m++
UC Irvine	http://www.ics.uci.edu/ mfukuda/m++

Table 3: M++ Web Pages

Download, unzip, and extract the file.

```
% gzip -d mpp.tar.gz | tar -xvf -
```

Once you have extracted the M++ archive or if you are a user at a site where M++ has been installed, you will find the *M++* directory that is organized as follows:

Directories	Content
m++/bin/	M++ executables (in the near future)
m++/src/	M++ daemon source and executables
m++/src/translator/	M++ compiler source and executables
m++/src/GMW1.xx/	GM zero-copy library (GMW1.11 for Solaris and 1.20 for Linux)
m++/sample/	Sample code
m++/appl/	Applications
m++/doc/	All manuals and papers

Table 4: Directory information

All executable codes will be generated in *m++/src* and *m++/src/translator* after installing M++.

2.2.2 Compiling gcc-2.95

The current implementation of M++ system is based on gcc-2.95.1 through to 2.95.1. If your site uses a newer version of gcc such as gcc-3.x, download gcc-2.95.3 and compile it.

1. Version Check
   Type **gcc -v** to see which version of gcc you are using. If it is newer than gcc-2.95.3, follow the instructions below, otherwise skip to Section 2.2.3.

2. Downloading
   Visit http://gcc.gnu.org/gcc-2.95/, download gcc-2.95.3 to the *m++* directory, and extract the files.

3. Compilation
   Assuming that your working directory include *m++* directory, compile the source code as instructed below:

```
cd m++/gcc-2.95.3
configure --prefix=/home/m++/gcc-2.95.3
make bootstrap
make install
cd ../..
```

2.2.3 Setting your shell environment

You need to set up your shell environment. Add the following two statements in your .cshrc and reinitialize it by invoking **"source /.cshrc"**.

Sites	A Statement to Add
medusa.bothell.washington.edu	set path=(/home/m++/gcc-2.95.3/bin $path /home/m++/src /home/m++/src/translator) setenv MPPDIR /home/m++/src setenv MPPARCH LINUX
howdy00.padc.mmpc.is.tsukuba.ac.jp	set path=(/home/m++/gcc-2.95.3/bin $path /home/mplus2/src /home/mplus2/src/translator) setenv MPPDIR /home/mplus2/src setenv MPPARCH SOLARIS_7
**.ics.uci.edu	set path=(/home/m++/gcc-2.95.3/bin $path /extra/mfukuda0/m++/src /extra/mfukuda0/m++/src/translator) setenv MPPDIR /extra/mfukuda0/m++/src setenv MPPARCH SOLARIS_2_5
**.cs.ehime-u.ac.jp	set path=(/home/mes/gcc-2.95.3/bin $path /home/mes/src /home/mes/src/translator) setenv MPPDIR /home/mes/src setenv MPPARCH SOLARIS_2_5
Other Linux systems	set path=(.../m++/gcc-2.95.3/bin $path .../m++/src .../m++/src/translator) setenv MPPDIR .../m++/src setenv MPPARCH LINUX

Table 5: Path set-up

For your information, a shell variable *$MPPARCH* specifies the underlying OS type as shown below:

Values	Machine Architecture
SOLARIS_7	PCs under Solaris7 or Solaris8
SOLARIS_2_5	Sparc and AP3000
LINUX	PCs under Linux

Table 6: $MPPARCH

If you are a user at any of the sites: UWB, Tsukuba, UCI and Ehime, skip to Section 2.4, otherwise go on to the next section to complete your M++ installation.

2.2.4 Compiling M++

Before compiling M++, you need to modify *Makefile* located in the *m++/src* directory. This file include the following commented-out statements:

```
#SOLARIS_2_5
#LIBS    = -lnsl -lsocket -lposix4 -ldl -lpthread -lfl
#SOLARIS_7
#LIBS    = -lnsl -lsocket -lposix4 -ldl -lrt -lpthread -lfl
#LINUX
#LIBS    = -lnsl -ldl -lrt -lpthread -lfl
```

Each statement corresponds to a different operating system type. Choose the appropriate one and remove the # character from the line head, so that this line will get effective. If your system is PC Linux, you should choose:

```
#LINUX
LIBS    = -lnsl -ldl -lrt -lpthread -lfl
```

Now, assuming that your working directory includes the *m++* directory in it, compile the source code as instructed below:

```
% cd m++/src
% make
% cd translator
% make all
% cd ../../..
```

2.2.5 Using GM library

If your cluster has installed a Myrinet card to each computing node, you can compile M++ with the underlying GM library for better performance. If you wish to use the GM library, follow the instruction shown below, otherwise skip to Section 2.4.

1. Obtaining the GM library
   Visit the Myricom homepage, *http://www.myri.com* and download the appropriate library. Note that the library downloading asks you to type a password received from Myrinet upon your purchase of Myrinet.

2. Installing the library
   Type the following commands to extract the GM library files in the GM home directory on your system:

   ```
   gunzip -c gm-1.5.1_Linux.tar.gz | tar xvf -
   cd {GM_HOME}
   ```

   Note that gzipped file name depends on the OS version The above file is valid for Linux. Thereafter, follow the *README-linux* file for the installation.

3. Modifying *m++/src/Makefile*

GM_LIBS	= -L{your GM library's directory path}/libgm/
GM_INCLUDES	= -I{your GM library's directory path}/include/
GMW_LIB	= -L{your M++ directory path}/src/GMW1.20/lib/
GMW_INCLUDE	= -I{your M++ directory path}/src/GMW1.20/lib_include/

   For instance, medusa.bothell.washington.edu has set up this *Makefile* as follows:

   ```
   GM_LIBS     = -L/usr/src/gm-1.5.1.1_Linux/libgm/
   GM_INCLUDES = -I/usr/src/gm-1.5.1.1_Linux/include/
   GMW_LIB     = -L/home/m++/src/GMW1.20/lib/
   GMW_INCLUDE = -I/home/m++/src/GMW1.20/lib_include/
   ```

4. Modifying *m++/src/GM1.20/Makefile*

GM_LIBS	= -L{your GM library's directory path}/libgm
GM_INCLUDES	= -I{your GM library's directory path}/include

   For instance, medusa.bothell.washington.edu has set up this *Makefile* as follows:

   ```
   GM_LIBS     = -L/usr/src/gm-1.5.1.1_Linux/libgm
   GM_INCLUDES = -I/usr/src/gm-1.5.1.1_Linux/include
   ```

5. Compiling the M++ source code

   Assuming that your working directory includes the *m++* directory in it, type the follow sequence of commands for compilation. After the compilation, the *m++/src* directory must include *thrd_gmw* and *inject_gmw*.

```
% cd m++/src/GMW1.20
% make
% cd ..
% make gmw
% cd ../..
```

2.3 Porting M++

Before porting M++, follow the instructions for system installation described in Section 2.2 to download and extract the tar file. Here are several tips to port M++ to your machine.

Hardware-dependent code

- setjmp() and longjmp() to capture and resume register values
- the M++ original thread library named *sthread*

Unix-dependent code

- dynamic-linking library, Pthread, and socket communication.

Among them, what you have to pay your attention for includes *setjmp()/longjmp()* functions, *sthread* and *pthread*.

The *setjmp()/longjmp()* pairs can be found in the files: SelfMigThr.h, thrFunc.h, and Sthread.cpp. Register contents are saved in a *jmpbuf* data structure by *setjmp* and resumed by *longjmp*. You need to know which register content will be saved in which *jmpbuf* element. Also, you must check how a stack grows. With these knowledge, modify those three cpp files.

The latter can be located in many files, however of importance is port.cpp. This code handles multiple sockets concurrently using *pthread*. Check the scheduling policy of your *pthread* library. If the library does not always force each thread to relinquish CPU whenever it is blocked on an I/O operations, the system will hang up. In that case, modify port.cpp so that no *pthread* will keep holding CPU upon an blocked I/O.

2.4 Additional Notes

1. If you want to run sample programs and/or applications, you are also required to visit each directory and to recompile the source code. Follow Section 3.2.
2. If you are a user at any of specific sites: UW Bothell, Univ. Tsukuba, UC Irvine and Ehime Univ., the following platform information may be useful.

(a) **Specification of UW Bothell's Medusa Cluster**

Spec. items	Descriptions
Node	PC (dual-i686 1.4GHz with 256KB cache/512MB main memory)
♯nodes	6 (medusa, mnode1 - mnode5)
Interconnection	4Gbps Myrinet [BCF+95]
File	NFS mounted to medusa's 34GB file system
External connection	medusa only

(b) **Specification of Tsukuba Univ.'s Howdy Cluster**

Spec. items	Descriptions
Node	PC (Athlon 1Ghz with DDR-SDRAM 256M-byte memory)
♯nodes	8 (howdy00 - howdy07)
Interconnection	1.28Gbps Myrinet [BCF+95]
File	NFS mounted to howdy00's 6.4GB file system
External connection	howdy00 only

(c) **Specification of UW Irvine's Iris Cluster**

Spec. items	Descriptions
Node	Sun Sparc Station SS5 with 96M-byte memory
♯nodes	9 (iris00 - iris09)
Interconnection	100Mbps Ethernet
File	NFS mounted to hippocrates' 1.6GB file system
External connection	Any nodes available from other ICS workstations

(d) **Specification of Ehime Univ.'s AP3000 Cluster**

Spec. items	Descriptions
Node	UltraSPARC 167MHz with X-byte memory
♯nodes	25 (apnet0000 - apnet0024)
Interconnection	1.6Gbps AP-Net
File	NFS mounted
External connection	Nodes 0, 1, 2, 3, 4, 8, 9, 10, 11, 16, 17, 18, and 19

3 Getting Started

3.1 M++ Profile

Upon initialization, an M++ daemon reads a configuration file. This file must be located as *profile* in each user's current working directory, (where the daemon is invoked.) It declares the following items:

Items	Descriptions
port	A socket port number used for thread migration.
	No default port. It must be given by a user.
hostnum	The number of PCs constituting the system.
	It must be 1 or larger.
mode	A thread library name: pthread or sthread.
	Pthreads are default. *Sthreads* are our original threads.
ip name and daemon ID	Each PC's IP name and its corresponding daemon ID.
	The daemon ID must be a unique non-negative integer.

The following *profile* defines a configuration in that an M++ daemon runs on howdy00, 01, 02, and 03, communicates with the other daemons through port# 92767, and simulates your agents with *Pthreads*.

```
port 92767
hostnum  4
mode  pthread
describe daemonid
howdy00  0
howdy01  1
howdy02  2
howdy03  3
describe_end
```

3.2 M++ Thread Compilation

Prior to an injection, each M++ self-migrating thread must be compiled into executable code with the *M++* compiler. The M++ compiler actually translates your M++ code into C++ programs and thereafter passes them to the g++ compiler. It accepts the following eight translation options:

options	remarks
-translate	translate M++ code to the corresponding C++ code and save it in files
-I	specify an *include* directory
-L	specify a *library* directory
-l	specify a *library* file
-D	make a given macro effective
-U	make a given macro ineffective
-E	preprocess M++ code and print out the result
-v	print out the version/translation information while compiling M++ code into the executable

For example, assume that your thread is coded in the *mythread.mpp* source file in the current working directory. If it refers to the math library, you can compile *mythread.mpp* to the *mythread* executable code by typing:

```
% m++ mthread.mpp -lm
```

Currently, the *M++* compiler has two restrictions:

1. A *.mpp* source code takes a different keyword to incorporate its header files. It is the *import*{ } keyword.

   - *import "Filename"* { *class/struct/union prototype declarations* } incorporates class/struct/union data structures from a user-defined *Filename.* If you incorporate only functions and macros, leave blank between this pair of blankets { }. However, note that { } must not be omitted.
   - *import <Filename>* { *class/struct/union prototype declarations* } incorporates class/struct/union data structures from a system-provided *Filename.* The notation, *import <Filename>* { } incorporates only the functions and macros of *Filename.*

   ```
   Example:
       import <cmath> {}
       import <cstdlib> {}
       import "Timer.h" { class Timer; }
   ```

2. The following keywords are reserved:

along	asm	auto	bool	break	case	catch
char	class	clocking	const	const_cast	continue	create
daemon	daemonobj	default	delete	destroy	do	double
dynamic_cast	edestroy	efork	eforkalong	ehop	ehopalong	ecreate
else	enum	explicit	extern	false	float	for
fork	forkalong	friend	goto	gvt_delta	gvt_end	gvt_ready
gvt_start	gvt_time	hop	hopalong	if	import	inline
int	link	long	mpp	mutable	namespace	new
node	operator	private	protected	public	register	reinterpret_cast
return	short	signed	sizeof	sleep	static	static_cast
struct	switch	template	terminate	this	thread	thread_id
throw	to	try	true	typedef	typeid	typename
union	unsigned	using	virtual	void	volatile	wakeup
wakeupall	wchar_t	while	with			

For example, consider compilation of *m++/sample/Car/Car.mpp.* This file includes three types of M++ threads:

Thread names	Descriptions
GSCashier	work as a gas station cashier who pumps gas into a tank
SportsCar	function as a sports car that gets regular fuel from a GSCashier
Truck	function as a truck that receives diesel fuel from a GSCashier

Compiling *m++/sample/Car/Car.mpp* is very simple. Copy the file to your current working directory and just type **"m++ Car.mpp"**. The compiler will produce three independent executables: *GSCashier*, *SportsCar*, and *Truck*, which are then ready to be injected to one of active M++ daemons.

3.3 System Invocation/Termination

Invocation

You must manually invoke an M++ daemon from all computing nodes that are to run your application. For each machine, follow the invocation procedure shown below:

1. Open an *xterm* or *tera-term* window.
2. Login the gateway or the master node of your cluster with ssh, (e.g., *medusa.bothell.washington.edu at UW Bothell*).
3. Login each machine with rsh.
4. Type *thrd* (or *thrd* & for a background execution).

The M++ daemon works as a simple Unix user process. It runs without being idle whether or not it has received active M++ threads to execute. If you debug your M++ threads or must share computing nodes with other users, you may want to start a daemon with a lower priority through a Unix *nice* option:

```
% nice 20 thrd&
```

When all the participating computing nodes have started a daemon, these daemons synchronize with one another and thereafter display the following message. (The example shows the message when four daemons are invoke.)

```
% thrd
port : 92767
Connection established to howdy00 with daemon 0
Connection: Wait for howdy01 acceptance ...
```

```
    Connection established to howdy01 with daemon 1
    Connection: Wait for howdy02 acceptance ...
    Connection established to howdy02 with daemon 2
    Connection: Wait for howdy03 acceptance ...
    Connection established to howdy03 with daemon 3
    Total daemons : 4
    Mode : pthread
    Stack : 65536 bytes
    Ready for being injected.
```

At this moment, you are ready to inject M++ threads.

Injection

The *inject* command injects new M++ threads to the system. The input format of this command is:

```
    % inject hostname filename #threads arg1 arg2 arg3 ....
```

The *hostname* parameter is the IP name of a target daemon where you want to inject M++ threads. The *filename* parameter specifies the thread executable code. The ♯*threads* parameter gives the number of threads to instantiate from *filename* and to inject into that daemon. Arguments such as *arg1*, *arg2*, and *arg3* are passed to each of those threads in form of a string array. This argument-passing mechanism is detailed in Section 4.

The following example injects M++ threads into *howdy00* to run the Car sample program.

```
    mplus2@howdy00[1]% thrd&
    [1] 18170
    port : 10145
    Connection established to howdy00 with daemon 0
    Total daemons : 1
    Mode : sthread
    Stack : 32768 byte
    Ready for being injected.

    mplus2@howdy00[2]% inject howdy00 ./GSCashier 1
    Injected 1 ./GSCashier thread(s) to howdy00
    GSCashier: Waiting for cars.

    mplus2@howdy00[3]% inject howdy00 ./Truck 1
```

```
Injected 1 ./Truck thread(s) to howdy00
Truck: I'm at the GasStation.
GSCashier: Injected fuel to Truck
Diesel: 900  Regular: 1000
GSCashier: Waiting for cars.
Truck: I'm filled with fuel.

mplus2@howdy00[4]% inject howdy00 ./SportsCar 2
Injected 2 ./SportsCar thread(s) to howdy00
SportsCar: I'm at the GasStation.
GSCashier: Injected fuel to SportsCar
Diesel: 900  Regular: 950
GSCashier: Waiting for cars.
SportsCar: I'm filled with fuel.

SportsCar: I'm at the GasStation.
GSCashier: Injected fuel to SportsCar
Diesel: 900  Regular: 900
GSCashier: Waiting for cars.
SportsCar: I'm filled with fuel.

mplus2@howdy00[5]% fg
~naoya/m++/thrd
^Cthrd: Interrupted.
thrd: Closed all connections.
thrd: Terminated.

mplus2@howdy00[6]%
```

This example assumes that the current directory has the three executable files of the Car sample program: *GSCashier*, *SportsCar*, and *Truck*. A user first invokes a single M++ daemon on howdy00. Second, he/she injects a *GSCashier* thread that behaves as a gas station cashier, initially prepares 1000 litters of diesel and regular gasoline respectively, and pumps gasoline into a car tank. Third, the user injects a *Truck* thread that receives 100 litters of diesel fuel from the *GSCashier* thread. Fourth, he/she injects two *SportsCar* threads, each receiving 50 litters of regular gasoline.

The *inject* command displays the following messages when detecting errors in thread injection:

- *SharedLib: Cannot load the dynamic module:* means that a given file name does not include executable thread code.

- *Segmentation fault:* means that the number of arguments actually given is incompatible to that of arguments coded in the thread program.

They are not fatal errors to M++ daemons. All you need is simply retype a correct file name and arguments.

<u>**Termination**</u>

One daemon termination is enough to kill all the other daemon processes. To do that, you must select an xterm/tera-term window where one of active daemons is running, bring the daemon in a foreground execution by the Unix **"fg"** command if it is running in the background, and terminate it by typing **"control + c"**. This interrupted daemon immediately forwards the signal to all the other daemons to terminate themselves. Upon a normal termination, the daemon displays the following message:

```
^Cthrd: Interrupted.
thrd: Closed all connections.
thrd: Terminated.
```

An unexpected trap/interrupt or the Unix *kill* command may cause an abnormal termination and keep the previous socket connection in use. This may also cause several daemons to remain alive. When it happens accidentally, you are required to:

1. visit all the xterm windows where those daemons are still running
2. kill those daemons by typing **"control + c"**, and thereafter
3. wait for a few minutes until Unix recognizes that the previous socket port is already released, or edit your *profile* to change the port number.

3.4 Using Myrinet Zero-Copy Communication

The M++ daemon communicates with others using either the TCP/IP or the Myrinet GM library. The latter implements zero-copy communication in that so-called *pin-downed* pages are allocated independently of ordinary OS-managed pages, and are accessed directly by an underlying Myrinet network interface card. The M++ daemon has an option to schedule its threads in such *pin-downed* pages that the network interface card can transfer those threads without OS interventions. We have confirmed in [FSCK01] that this scheme improved migration performance 1.85 times better than using TCP/IP in a multi-agent application.

To implement inter-daemon communications with the GM library, use the commands listed below:

Procedures	Corresponding
Daemon compilation	make gmw
Daemon invocation	thrd_gmw
Thread injection	inject_gmw hostname filename #threads arg1 arg2 arg3 ....

4 Writing M++ Threads

This section describes how to program M++ threads. You will first capture the overview of M++ executable model, thereafter understand how to define logical network components, and finally learn how to program M++ thread code.

Before going on to the following sub-sections, please be reminded of the usage of a pair of square brackets [ ] or *[ ]*:

- a pair of square brackets in italic *[ ]*
  Symbols inserted between a pair of italic brackets *[ ]* may be omitted.
- a pair of square brackets in bold **[ ]**
  this pair of brackets is a part of M++ statement or method.

4.1 Execution Model

M++ involves three levels of networks as shown in Figure 1. The lowest level is the *physical network* (Myrinet in both Tsukuba's *howdy* and UWB's *medusa* PC clusters, AP-NET in Ehime's AP3000, and 100Mbps Ethernet in UCI's *iris* Sparc cluster), which constitutes the underlying computational nodes. Superimposed on the physical layer is the *daemon network*, where each daemon is a Unix process executing and exchanging M++ threads with others. As described in Section 3.1, its processor mapping and system unique ID, termed *daemon ID*, are predefined in a user profile. The *logical network* is an application-specific computation network dynamically constructed by M++ threads on top of the daemon network. Each logical node has a *node ID* local to the corresponding daemon, while a logical link maintains a set of *source* and *destination IDs*, each local to the current and the destination logical node respectively.

M++ threads distinguish four C++ classes of network objects: *daemon*, *daemonobj*, *node*, and *link*. The *daemon* object provides M++ threads with the current daemon status. The *daemonobj* object stores user-defined data shared by all M++ threads

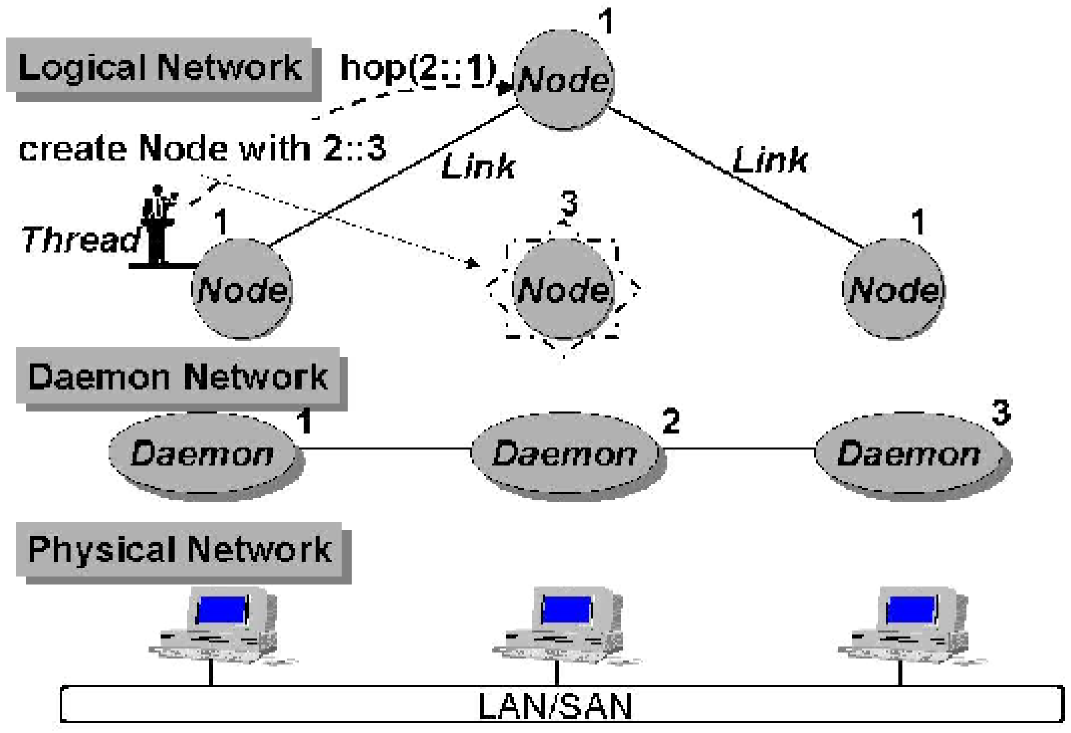

Figure 1: M++ network architecture

running on the same daemon process. The *node* object represents a logical network node, whose method and data members are accessed by M++ threads residing on this node. The *link* corresponds to a logical network link, makes its method and data members visible to threads residing on the both ends, (i.e., nodes) of this link. The *daemon*, *node*, and *link* objects also provide M++ threads with system-defined methods to inform of their network status: the current daemon ID, the total number of daemons, the host name, the node ID, the node class name, the number of threads on the current node, etc.. Details are given in Section 4.2. Sharing with these network objects is considered as a course of inter-thread communication, which however does not mean that M++ threads are provided with a complete view of distributed shared memory. Prior to accessing a given object, they must migrate themselves to it.

Figure 2 shows the M++ thread framework. A user can define various threads using the *thread* keyword. As in an ordinary C++ class, each thread allows its data members and member functions to be read and written in three access modes such as *public*, *protected*, and *private*. Their default attribute is *private*. The thread carries all data members with it whenever moving to a different node. However, it does not carry its member functions, which means that the same executable file must be made available by NFS or manually copied in advance at at each destination. The thread does not carry any local variables, either. This means that, while local variables including arguments can be still defined inside each member function, the validity of their contents is guaranteed only till a next thread migration.

The thread has constructors which can accept arguments. When a constructor is invoked with two arguments such as *(int)argc* and *(char **)argv* as shown in the following code, it can handle arguments passed from the *inject* command. In that case, the meaningful arguments start from *argv*[*4*]. (In precise, *argv*[*0*], *argv*[*1*], *argv*[*2*], and *argv*[*3*] include a string *"inject"*, *a target host name*, *this thread executable file*, and *# threads to be injected* respectively.) **The destructor has not yet been supported by the current version of M++. It will be made available in the following version.**

The *void main( )* function is the core of the thread, describing its autonomous behavior involved in network construction, navigation, inter-thread communication, and computation. It is automatically called by the system immediately after a constructor invocation. In Figure 2, a thread hops to a system-predefined node *INIT* on the daemon #0, and calls its *work* function. A return from the *main( )* means the termination of the corresponding thread. Multiple inheritance and polymorphisms are available in the M++ thread. However, one exception is that *main( )* cannot be inherited from base threads, (in other words, a sub-thread must define its own *main( )*.)

```
thread SimpleThread {
public:
  SimpleThread( int argc, const char** argv ) : num( atoi( argv[4] ) )
                { }
  void main( ) { hop( INIT @ 0 );
                 work( );
               }
private:
  void work( ) { cout << num << endl;
               }
  int num;
};
```

Figure 2: Thread Framework

Upon a startup, the system creates a logical node named *INIT* on every daemon. Any M++ thread may be injected (by the *inject* command from the shell or by another M++ thread) into any of the *INIT* nodes. It may then start network construction and migration using the following five groups of methods: *[e]create*, *[e]destroy*, *[e]hop*, *[e]fork*, and *[e]inject*. Note that the methods starting from 'e' return an error code, otherwise they return nothing for performance reasons.

- *[e]create node( );* permits an M++ thread to create a new logical node. See Section 4.2.3 for more details.
- *[e]create link( );* permits an M++ thread to create a new logical link. See Section 4.2.4 for more details.
- *[e]destroy node( );* permits an M++ thread to destroy an existing logical node. See Section 4.2.3 for more details.
- *[e]destroy link( );* permits an M++ thread to destroy an existing logical link.See Section 4.2.4 for more details.
- *[e]hop( );* navigates an M++ thread over an existing logical network. Its execution resumes from the statement following this hop keyword. See Section 4.3.6 for more details.
- *[e]fork( );* allows an M++ thread to spawn its copy on a specified destination and have it start from the statement following the fork method. The original thread continues its execution without waiting for the duplicated thread to be terminated. See Section 4.3.6 for more details.
- *[e]create thread( );* starts a different M++ thread from the beginning of its code on a specified destination. The original thread continues its execution without waiting for the injected thread to be terminated. See Section 4.3.2 for more details.

4.2 Network Components

This section gives the procedures to define and access the four M++ classes of network objects: *daemon*, *daemonobj*, *node*, and *link*. Note that, if those classes are defined in a header file, you must incorporate them into your *mpp* source code by enumerating all those class names in a *import* keyword. See Section 3.2 for the complete explanation.

4.2.1 The *daemon* Class

M++ threads can obtain the current daemon status at any point of time by accessing the system-provided *daemon* class object. This class includes the following public functions that can be accessed through the *daemon* keyword followed by a dot "." and the corresponding function name.

- *string name( );* returns the IP name of the current daemon.
- *int id( );* returns the current daemon ID.

- *int total( );* returns the total number of daemons in the system.

4.2.2 The *daemonobj* Class

Unlike the system-provided *daemon* class, this is a user-provided class of which each M++ daemon instantiates only one object upon receiving the first M++ thread that has defined this *daemonobj* class. Thereafter, all M++ threads with the same *daemonobj* class definition can access this *daemonobj* instance as far as they reside on the same daemon. A user can define as many distinct *daemonobj* classes as needed, using the ordinary C++ class definition. However, each different M++ thread must declare which *daemonobj* class to use *a priori*. Such a declaration is achieved using the *daemonobj* keyword followed by a *daemonobj* class constructor call in front of the M++ thread definition. This is where the constructor receives its arguments which must be however literal values only.

Note that the classes defined for *daemonobj*s can be derived using single/multiple inheritance and from abstract classes as in C++. Anyway types of data members, method arguments, and pointers are permitted unless they references M++ threads which may migrate somewhere.

```
class A {  // The daemonobj A is declared as an ordinary C++ object.
public:
  A( int a, string b ) : { }
  ....;
};
daemonobj A( 0, ‘‘abc’’ ); // A is the daemonobj which the thread T uses.
thread T {
  ....;
};
```

The *daemonobj* keyword is also used as the reference name of the *daemonobj* object. No matter what class a thread has declared as its *daemonobj*, it must use the reference name *daemonobj* in order to access the object. There are no other reference names than *daemonobj*. The following code example shows two different *daemonobj* class definitions such as *DaemonObj1* and *DaemonObj2*, each respectively used by the *Thread1* and *Thread2* M++ threads:

```
class DaemonObj1 {
public:
  DaemonObj1( int num ) : num( num ) { }
  int num;
};
```

```
class DaemonObj2 {
public:
  DaemonObj1( const char* str ) { strcpy( DaemonObj2::str, str ); }
  char str[20];
};

daemonobj DaemonObj1( 0 );
thread Thread1 {
public:
  void main( ) { cout << ``I'm using DaemonObj1: '' << ++daemonobj.num << endl; }
};

daemonobj DaemonObj2( ``Hello'' );
thread Thread2 {
public:
  void main( ) { cout << ``I'm using DaemonObj2: '' << daemonobj.str << endl; }
};
```

When receiving the first *Thread1* thread, a daemon instantiates a *DaemonObj1* object by invoking its constructor with 0 and resumes the *Thread1* thread that thereafter increments the *num* variable to 1. The daemon however will no longer instantiate a new *DaemonObj1* object but simply passes this reference to the second or the following *Thread1* threads that visit this daemon. Similarly, the daemon instantiates a *DaemonObj2* object with a string constant, "hello" when receiving the first *Thread2* thread. Thus, all the *Thread2* thread visiting this daemon retrieves "hello" as the content of *daemonobj.str*.

Note that each *daemonobj* objects, once allocated to a daemon, will never been deallocated until the system is terminated.

4.2.3 Nodes

The M++ *node* represents an application-specific logical network node. It is a place where M++ threads migrate themselves and carry out their task. A user can define various C++ classes for and instantiate them as M++ *node*s. However, their actual instantiation takes place in response to runtime requests ordered by M++ threads. (This in turn means that no *node*s are instantiated at compile time.) Once a *node* is instantiated from a user-defined C++ class, M++ threads may migrate to this *node*. All threads residing on the same *node* can share its data and method members. They are guaranteed to access the current *node* without intervened by any other threads unless they call either a migration or a synchronization method (see Section 4.3.7).

The following code shows an example of a class defined for *node*s. Similar to classes used for *daemonobj*, it can be derived using single/multiple inheritance and from abstract classes as in C++. Any types of variables and arguments are allowed except pointers referencing M++ threads.

```
class Place {
public:
  Place( int num ) : num( num ) { }
  void setNum( int num ) { Place::num = num; }
  int  getNum( )         { return num; }
  void printName( )      { cout << ``I'm `` << name( ) << endl; }
private:
  int num;
}
```

Note that this *Place* class includes the *name( )* method whose body is not actually specified inside this class definition. It is one of the system-predefined methods that are available to any classes used for M++ *node*s:

- *int id( );* returns the current node ID.
- *string name( );* returns the class name from which the current node was instantiated.
- *int thr_num( );* returns the total number of threads residing on the current node.

M++ threads are in charge of choosing which class it uses for a new *node* instantiation, instantiating an object from the class, and mapping it onto a specific daemon as well as giving a *node ID* to this object. To perform such a node creation, M++ threads must call the following method in its main( ):

- *[e]create node<NodeClassName>(args_list) with (NodeID[@DaemonID]);*

where *NodeClassName* is the name of the class used for a *node* instantiation; a*args_list* is a list of arguments passed to this class constructor; *NodeID* is a daemon-unique node identifier that must be a non-negative integer; and *DaemonID* is the identifier of the daemon to which this node is mapped. If *DaemonID* is omitted, a new node is created on the current daemon. This method returns *mpp::Errno::noError* on success, otherwise *mpp::Errno::createNodeFailed.* If *e* is not prefixed to *create*, the return value is always *mpp:Errno::noError.*

In order to access a *node*, M++ threads must move to the node by calling one of its migration methods. One of such methods is:

- *[e]hop([NodeID[@DaemonID]);*

where *NodeID* and *DaemonID* are the destination node's and the destination daemon's identifier respectively. Again, the omission of *DaemonID* specifies the current daemon as a default. This method returns *mpp::Errno::noError on success*, otherwise *mpp::Errno::hopFailed*. It always returns *mpp::Errno:noError* if it has no *e* prefix. To understand all thread migration methods in details, see Section 4.3.6.

M++ threads can access their current *node* through the reference name, *node*. There are no reference names other than *node*

- *node.member*: where *member* is one of system-provided node methods such as *id( )*, *name( )*, and *thr_num( )*.
- *node<NodeClass>.member*: where *NodeClass* is the class name from which this node was instantiated, and *member* is a member of this class.

M++ threads are also responsible to delete existing *node*s. A *node* deletion is performed by:

- *[e]destroy node(NodeID[@DaemonID]);*

This method deletes a node with *NodeID* mapped on the daemon with *DaemonID*. (The omission of *DaemonID* means the current daemon.) It returns *mpp::Errno::noError* on success or always *mpp::Errno::noError* without the 'e' prefix, otherwise *mpp::Errno::deleteNodeFailed*. If 'e' is not prefixed, the return value is always *mpp::Errno::noError*. Note that if a node to be deleted has threads and *link*s, it cannot be deleted. This also means that an M++ thread cannot delete the current *node*.

The following code example shows a sequence of node creation, access, and deletion:

```
class Place {
public:
  int num;
};

thread SimSpace {
public:
  void main( );
private:
  int i;
}

void SimSpace::main( ) {
  create node<Place> with( 0 );
  hop( 0 );
  node<Place>.num = daemon.id( );
  hop( INIT );
  destroy node( 0 );
}
```

The *SimSpace* thread first creates a new *node* from the *Place*, gives the nodeID '0' to this node, and maps it onto the current daemon. Thereafter, the thread migrates to this node and substitutes this node's *num* variable with the current daemon id which is retrieved by *daemon.id( )*. Finally it hops back to the system-predefined *INIT* node and deletes the node whose id is 0.

4.2.4 Links

The M++ *link* represents a logical path along which M++ threads move from one *node* to another. While M++ threads can jump directly to a *node* addressed with its absolute node/daemon identifiers, *link*s are another addressing option that permits threads to move to a *node* neighboring to the current *node*, (e.g., move to a left neighbor or a upward neighbor). Similar to the *node*, a user can define various C++ classes for and instantiate them as M++ *link*s. Again, like *nodes*s, the actual instantiation of *link*s must be requested from and achieved by M++ threads at run time. Once a *link* is instantiated from a user-defined C++ class, M++ threads may migrate along this *link*. A *link* is accessible from all threads residing on either one of the two end points of the same *link*, (i.e. one of the two *node*s incident to the same *link*). Note that the system guarantees a non-interruptible access to a *link*. In other words, while a thread is executing a *link* method, no other threads can access any methods of the same *link*.

Each *link* has two identifies: source link ID and destination link ID. The former is unique to the source node, the latter unique to the destination node. Of notable is that the source and the destination IDs are recognized as the destination and the source IDs at the destination node. The link ID must be a non-negative integer.

The following code shows an example of a class defined for *link*s. The C++ class inheritance is possible. Any types of variables and arguments are allowed except pointers referencing M++ threads.

```
  class Path {
  public:
    Path( int num ) : num( num ) { }
    void setNum( int num ) { Place::num = num; }
    int  getNum( )         { return num; }
    void printName( )      { cout << ``I'm ''    << name( )    << endl; }
    void printId( )        { cout << ``src=''    << src_id( )
```

```
                                    << ‘‘ dest= ’’ << dest_id( ) << endl; }
  private:
    int num;
  }
```

Note that this *Path* class includes *name( )*, *src_id( )*, and *dest_id( )* whose body is not actually specified inside the class definition but implicitly provided by the system. Two groups of methods have been made available to M++ *link*s: *link-shared* and *link-local* groups. The former provides general information shared among all *link*s emanating from the same *node*, while the latter focuses on the information local to each *link*.

The following list summarizes **link-shared** methods:

- *int count( );* returns the total number of links emanating from the current node.
- *int min( );* returns the minimal link ID.
- *int max( );* returns the maximal link ID.
- *int next(int linkId);* returns the minimal link ID among those which are larger than *linkId*.
- *int prev(int linkId);* returns the maximal link ID among those which are smaller than *linkId*.
- *bool exists(int linkId);* returns true if there exists the link with *linkId*.

The following list summarizes **link-local** methods:

- *int name( );* returns the class name from which the current *link* was instantiated from.
- *int src_id( );* returns the source link ID.
- *int dest_id( );* returns the destination link ID.
- *int src_daemon_id( );* returns the source daemon ID.
- *int dest_daemon_id( );* returns the destination daemon ID.
- *int src_node_id( );* returns the source node ID.
- *int dest_node_id( );* returns the destination node ID.

M++ threads are in charge of choosing which class it uses for a new *link* instantiation, instantiating an object from the class, and mapping it onto a specific daemon as well as giving a pair of source and destination *link ID*s to this object. To perform such a link creation, M++ threads must call the following method in its main( ):

- *[e]create link<LinkClassName>(args_list) with (SrcID) to (NodeID[@DaemonID]) with (DestID);*

where *LinkCalssName* is the name of the class used for a *link* instantiation; *args_list* is a list of arguments passed to this class constructor; *SrcID* is a link ID unique to the source node; *NodeID* is the destination node identifier; *DaemonID* is the destination daemon's identifier; and *DestID* is a link ID unique to the destination node. The omission of *DaemonID* means that the destination node is on the current daemon. This method returns *mpp::Errno::noError* on success, otherwise *mpp::Errno::createLinkFailed*. If 'e' is not prefixed to *create*, the return value is always *mpp::Errno::noError*.

Note that the source and the destination node of a *link* may be identical. In other words, the system permits threads to create a circle *link* departing from and arriving at the same *node*, even in which case the source and destination IDs of this *link* must be distinguished.

Once a *link* has been created, M++ threads can move to another *node* along the *link* by calling one of its migration methods. One of such methods is:

- *[e]hopalong(srcLinkID);*

where *srcLinkID* is the identifier of the link which emanates from the current node and along which the calling thread will move to its destination. This method returns *mpp::Errno::noError* on success, otherwise *mpp::Errno::hopFailed*. It always returns *mpp::Errno:noError* if it has no 'e' prefix. To understand all thread migration methods in details, see Section 4.3.6.

All *link*s incident to the current *node* can be accessed through the reference name, *link*. There are no reference names other than *link*

- *link.member*: where *member* is one of link-shared methods which the system has predefined: *count( )*, *min( )*, *max( )*, *next( )*, *prev( )*, and *exists( )*.
- *link*[*LinkID*]*.member*: where *member* is one of link-local methods which the system has predefined: *name( )*, *src_id( )*, *dest_id( )*, *src_daemon_id( )*, *dest_daemon_id( )*, *src_node_id( )*, and *dest_node_id ( )*.
- *link*[*LinkID*]*<LinkClass>.member*: where *LinkClass* is the class name from which this link was instantiated, and *member* is a member of this class.

M++ threads are also responsible to delete existing *link*s. A *link* deletion is performed by:

- *[e]destroy link(LinkID);*

This method deletes a link whose source link ID is *LinkID*. It returns *mpp::Errno::noError* on success or always *mpp::Errno::noError* without the e prefix, otherwise *mpp::Errno::deleteLinkFailed*. If 'e' is not prefixed, the return value is always *mpp::Errno::noError*.

The following code example shows a sequence of link creation, access, and deletion:

```
class Place {                  // a class used for node instantiation

};

class Path {                   // a class used for link instantiation
public:
  Path( ) : count( 0 ) { }
  int count;
}

thread SimSpace {              // a thread instantiate a node and a link
public:
  void main( );
private:
  int i;
}

void SimSpace::main( ) {       // a SimSpace thread's main behavior
  create node<Place> with( 0 );
  create link<Path> with( 0 ) to ( 0 ) with ( 1 );
  link[0]<Path>.count++;
  hopalong( 0 );
  cout << ``link.count = '' << link[0]<Path>.count++ << endl;
  hopalong( 1 );
  cout << ``link.count = '' << link[0]<Path>.count++ << endl;
  destroy link( 0 );
}
```

The *SimSpace* thread first instantiates a new *node* from the *Place*, gives the nodeID '0' to this *node*, and maps it onto the current daemon. Thereafter, it establishes a new *link* from the current *node*, (i.e., an INIT *node* where injected) to the *node 0*. This *link* is accessible as *link 0* from the INIT *node* and as *link 1* from *node 0*. Thereafter, the thread increments this link's *count* variable and hops it along to *node 0*. Again, it increments this link's *count* variable and hops it along back to the INIT *node*. Finally, it deletes this link.

4.2.5 Classes Used for Network Components

It is possible to use the same class for definition of various network components such as *daemonobj*, *node*, and *link*, as far as it is an ordinary C++ class. However, if a class uses network-component-specific methods such as *name* and *src_id*, it must be used for representing only the network components that provide those specific methods at the system level. For instance, *DaemonObj1* and *DaemonObj2* in Section 4.2.2 are an ordinary C++ class which does not include any system-provided methods, and thus they can be also used as both *node*s and *link*s. On the other hand, *Place* in Section 4.2.3 includes a system-provided method, *name( )* which is available for *node*s and *link*s, and therefore it must be used for *node*s and *link*s but not for *daemonobj*s.

4.3 Threads

M++ threads are active computation entities, capable of constructing and migrating over a logical network. They are defined using the *thread* keyword rather than *class*, however their definition takes a form similar to that of C++ classes except a few restrictions. Different types of M++ threads can be programmed inside the same *m++* program.

Upon an instantiation, they are initialized with their constructor and automatically activated to execute their *main* method. Those steps are similar to Java threads: a constructor initialization and a *run* method invocation, however note that M++ threads run their *main* method automatically, while Java threads' *run* method must be triggered by the *start* method.

The *main* method is the heart of M++ threads, and can perform:

- initialization of *deamonobj*s,
- creation of *node*s, *link*s, and other *thread*s,
- access to all M++ objects such as *daemon*s, *daemonobj*s, *node*s, *link*s, and other *thread*s
- destruction of *node*s, *link*s and other *thread*s,
- migration to another *node*,
- forking itself to the current or another *node*,
- direct communication with another *thread*, and
- inter-threads synchronization

Each of those manipulations listed here is explained in the following subsections.

4.3.1 Thread Definition

The definition of M++ thread starts from the *thread* keyword rather than *class*. It follows all C++ class rules for single/multiple inheritance, polymorphisms, and abstract definition, however there is one exception: the *main( )* method defined inside each *thread* will not be inherited by any sub-threads. All sub-threads must define their own *main( )* method.

If a thread definition has one or more pure virtual functions, it is regarded as an abstract thread and does not have to include its *main( )* method. Nevertheless to say, no instantiation is possible from an abstract thread. It is used to derive sub-threads which no longer include pure virtual functions.

Data members declared in *thread* are carried with each individual thread upon a migration. However, there are three restrictions in the use of variables:

- Local variables are **not** carried with each thread. They can be still declared and used inside each method, however their contents are not guaranteed once a thread has migrated to another *node*.
- Arguments passed to any methods defined in *thread* are **not** carried with a thread, either. Their contents are not guaranteed once a thread has migrated to another *node*.
- Pointer contents are not guaranteed once a thread has migrated to a remote *node*, although they can be still defined as *thread* data members. This implies that memory allocated within *thread* using *new* must be freed with *delete* before a next migration. (Note that dynamic memory allocation works in any other classes, since they are immobile.)

The following code defines the *Base* thread from which the *Sub* thread is derived as a sub-thread.

```
thread Base {                // a super thread
public:
  Base( int num ) : num( num ) { }
protected:
  int getNum( ) { return( num ); }
  void setNum( int num ) { Base::num = num; }
  virtual void printNum( ) = 0;
private:
  int num;
};

thread Sub : public Base {   // a sub thread
public:
  Sub( int argc, const char** argv ) : Base( atoi( argv[4] ) ) { }
  void main( ) { printNum( ); }
protected:
  void printNum( ) { cout << ``Num = `` << getNum( ) << endl; }
};
```

The *Base* thread includes a pure virtual function, *printNum( )*, because of which it does not have to include its own *main( )* method. The *Sub* thread inherits *Base* and implements the *printNum( )* method. Since it have no more pure functions, it must define the *main( )* method which simply calls *printNum( )*. Its constructor converts the *argv*[*4*] string into an integer that is passed to the *Base* constructor. This integer value is stored in the *num* variable and printed out when the *main( )* calls *printNum( )*.

4.3.2 Thread Creation and Constructor

M++ threads are instantiated either (1) by running the *inject* command from the Unix *shell* or (2) executing the *create* statement from any other M++ thread.

1. *inject*:
   The *inject* command creates one or more thread instances at a time. It takes the following arguments:

   ```
   % inject hostname ThreadClassName #instances arg1 arg2 .....
   ```

   where *hostname* is the IP name of a cluster computing node onto which you will instantiate M++ threads; *ThreadClassName* is the name of thread class; *#threads* is the number of instances created from *ThreadClassName*; and *arg1 arg2 ...* are the actual arguments passed to a constructor of each instance. All instantiated threads start at the *INIT* node on the designated computing node.

2. *create*:

   The *create* statement creates only one thread instance at a time. It takes one of the following three forms:

   - *[e]create thread<ThreadClassName>(args_list) to(NodeID[@DaemonID]);*
     where *ThreadClassName* is the name of thread class; *args_list* is a list of char[] arguments passed to a constructor, (i.e., a list of character arrays); and *NodeID[@DaemonID]* specifies a node where a thread is instantiated. If *DaemonID* is omitted, the instance is placed at a designated node on the current daemon.

- *[e]create thread<ThreadClassName>(args_list) along(LinkID);*
  where *ThreadClassName* is the name of thread class; *args_list* is a list of char[] arguments passed to a constructor, (i.e., a list of character arrays); and *LinkID* specifies a link whose destination node is where a new instance is created.
- *[e]create thread<ThreadClassName>(args_list)*
  where *ThreadClassName* is the name of thread class, and *args_list* is a list of char[] arguments passed to a constructor, (i.e., a list of character arrays). This statement instantiates a thread on the current node.

If the *create* statement has a prefix 'e', it returns *mpp::Errno::noError* on success, otherwise *mpp::Errno::injectFailed*. Without the 'e' prefix, it always returns *mpp::Errno::noError*.

Constructors of M++ thread may or may not accept arguments. If a thread is designed to always derive sub-threads, its constructors may accept any types of arguments. If a thread is however designed to accept arguments given from either *inject* or *create* upon its instantiation, it must include a constructor which receives *int argc* and *char** argv* arguments. The *argc* argument contains four plus the number of arguments passed from either *inject* or *create*: *arg1 arg2 ...*. The *argv* arrays include:

argv	when invoked from inject	when invoked from create
argv[0]	a string "inject"	null
argv[1]	the target host name where injected	null
argv[2]	this thread class name	null
argv[3]	# of instances (represented in a string)	null
argv[4]	arg1	arg1
argv[5]	arg2	arg2
...	...	...

The following code defines two threads such as *Parent* and *Child*. *Parent* simply creates 10 *Child* instances on the current node, and passes two arguments to them. The first argument is "I am Thread No." for all *Child* threads, whereas the second argument is "0", "1", "2", ..., "9" to each of them. Each *Child* stores the first argument in the *msg* string variable, converts the second argument in the *id* integer variable, and prints out *msg* and *id* in its *main( )* method.

```
thread Child {
public:
  Child( int argc, const char** argv ) : id( atoi( argv[5] ) ) {
    strncpy( msg, argv[4], 16 );
  }
  void main( ) { cout << msg << id << endl; }
private:
  char msg[16];
  int id;
}

thread Parent {
public:
  void main( );
private:
  int i;
  char numstr[10];
};

void Inject::main( ) {
  for ( i = 0; i < 10; i++ ) {
    sprintf( idstr, ``%d'', i );
    create thread<PrintMsg>( ``I am Thread No.'', idstr );
  }
}
```

4.3.3 Thread Termination and Destructor

The M++ thread is terminated either implicitly by reaching the end of its *main( )* method or explicitly by executing the following statement.

- *terminate;*
  terminates the invoking thread.

The current version of M++ has not supported the thread destructor. Although you may define it, nothing would happen.

4.3.4 Thread Status

Each M++ thread includes several system-provided methods that indicate the thread status. They are categorized into two groups: *thread-shared* and *thread-local* groups. The former provides general information shared among all *threads* residing on the same *node*, while the latter focuses on the information local to each *thread*.

The following list summarizes *thread-shared* methods:

- *int count( );* returns the total number of threads residing on the current node.
- *int min( );* returns the minimal thread ID.

- *int max( );* returns the maximal thread ID.
- *int next(thread_id threadId);* returns the minimal thread ID among those which are larger than *threadId*.
- *int prev(thread_id threadId);* returns the maximal thread ID among those which are smaller than *threadId*.
- *bool exists(thread_id threadId);* returns true if there exists the thread with *threadId*.

The following list summarizes *thread-local* methods:

- *int name( );* returns the class name from which this *thread* was instantiated.
- *int src_id( );* returns the identifier of this *thread*.

All *thread*s residing on the same *node* can access one another through the reference name, *thread*. There are no reference names other than *thread*.

- *thread.member*: where *member* is one of thread-shared methods which the system has predefined: *count( )*, *min( )*, *max( )*, *next( )*, *prev( )*, and *exists( )*.
- *thread*[*ThreadID*]*.member*: where *member* is one of thread-local methods which the system has predefined: *name( )* and *id( )*.
- *thread*[*ThreadID*]<*ThreadClass*>*.member*: where *ThreadClass* is the class name from which this thread was instantiated, and *member* is a member of this class.

4.3.5 Network Construction

The construction of logical networks is viewed as a series of creating and deleting *node*s and *link*s. Threads are only entities that can instantiate those network components.

- *[e]create node<NodeClassName>(args_list) with (NodeID[@DaemonID]);*
  Details are given in Section 4.2.3.
- *[e]destroy node(NodeID[@DaemonID]);*
  Details are given in Section 4.2.3.
- *[e]create link<LinkClassName>(args_list) with (SrcID) to (NodeID[@DaemonID]) with (DestID);*
  Details are given in Section 4.2.4.

- *[e]destroy link(LinkID);*
  Details are given in Section 4.2.4.

The following code permits an M++ thread to create a *MyNode* node on each daemon and establish a *MyLink* link to each of those nodes from the current working node, (i.e., the *INIT* node where it has been injected).

```
class MyNode { };            // a class used to instantiate a node

class MyLink { };            // a class used to instantiate a link

thread StarNet {             // a thread creating nodes and links
private:
  int i;
  int nodeId, linkId;
public:
  StarNet( int argc, char** argv ) :
    nodeId( atoi( argv[4] ) ), linkId( atoi( argv[5] ) ) { }
  void main( ) {
    for ( i = 0; i < daemon.total( ); i++ ) {
      create node<MyNode>( ) with ( nodeId @ i );
      create link<MyLink>( ) with ( i ) to ( nodeId @ i ) with ( linkId );
    }
  }
};
```

In this example, the *StarNet* thread receives two arguments, each stored into its *nodeId* and *linkId* variables. It creates a new *node* on each of all daemons whose identifier ranges from 0 to *daemon.total( )*. Each *node* receives the content of *nodeId* as its identifier. Thus, all the *node* obtains the same identifier while each resides on a different daemon. As soon as the thread creates a new *node*, it establishes a new *link* to this *node* from the current node. The source link ID of each *link* is set to its destination daemon's identifier, while the destination link ID is set to the content of *linkId*.

4.3.6 Thread Deployment

An M++ thread can migrate itself or deploy its copy over an existing logical network by calling one of the following methods. Similar to the network construction, if those methods are prefixed with 'e' and an error has occurred, they returns an error code. Otherwise, they return *mpp::Errno::noError*, (i.e., no errors.) For all those methods, if the *DaemonId* argument is omitted, their migration or deployment is bound for the local daemon. If the *NodeId* argument is omitted together, their destination is the current working node. Note that it is not permissible to omit only *NodeId*.

- *[e]hop([NodeID[@DaemonID]]);*
  migrates the calling thread to the node specified with *NodeId* and *DaemonId*. If this method is not prefixed with 'e' and an error has occurred at an migration, the calling thread is automatically terminated as displaying a termination message. If it is prefixed with 'e' and an error has occurred at an migration, the calling thread remains at the same node as receiving an error code.

- *[e]hopalong(SrcLinkID);*
  migrates the calling thread along the link specified with *SrcLinkId*. If this method is not prefixed with 'e' and an error has occurred at an migration, the calling thread is automatically terminated as displaying a termination message. If it is prefixed with 'e' and an error has occurred, the calling thread remains at the same node as receiving an error code.

- *[e]fork([NodeID[@DaemonID]]);*
  deploys a child of the calling thread to the node specified with *nodeId* and *daemonId*. The child thread starts its execution from the statement next to *[e]fork( )*. The parent thread keeps running without waiting for its child.

- *[e]forkalong(SrcLinkID);*
  deploys a child of the calling thread along the link specified with *linkId*. The child thread starts its execution from the statement next to *[e]forkalong( )*. The parent thread keeps running without waiting for its child.

The following code example permits an M++ thread to visit each leaf node of a star network already constructed by another M++ thread whose code was shown in Section 4.3.5.

```
thread StarHopper {
private:
  int i;
  int srcNode, srcDmn;
public:
  StarHopper( int argc, char** argv ) { }
  void main( ) {
    srcNode = node.id( );                           // record the origin node id
    srcDmn  = daemon.id( );                         // record the origin daemon id
    for ( i = 0; i < daemon.total( ); i++) {        // for each daemon
      hopalong( i );                                // hop along the link to it
      cout << "hello! at " << daemon.id( ) << endl; // display the dest. daemon id
      hop( srcNode@srcDmn );                        // go back to the origin node
    }
  }
}
```

This example first memorizes the node and daemon identifiers of where a *StarHopper* thread has been injected. It thereafter repeats dispatching the thread along each different link emanating from the origin, labeled from 0 to *daemon.total( )* - 1, thus leading to each leaf daemon. Every time the thread visits on a different leaf node, it displays “hello!” and jumps back to the origin node.

4.3.7 Synchronization

The following statements are used for inter-threads synchronization.

- *sleep [node | daemon]*
  suspends the execution of the calling thread which will sleep until it is woken up by either a *wakeup* or a *wakeupall* statement.
- *wakeup ThreadID | node | daemon*
  wakes up the thread specified by *ThreadID* if it is given, wakes up one of the thread sleeping on the current node if the *node* parameter is given, or wakes up one of those sleeping on the current daemon if the *daemon* parameter is given. The woken-up thread will resume its execution after the point where it was suspended. A sleeping thread is chosen to wake up in a first-come-first-service order so that no starvation occurs.
- *wakeupall node | daemon*
  wakes up all threads sleeping on the current node or all those sleeping on the current daemon according to a parameter given to the *wakeupall* statement. All those threads will resume their execution after the point where they were suspended.

Non-interruptible execution is guaranteed by implementing every node as a monitor. Each node allows only one of the arriving threads to resume its execution and access the node variables. Thus, there is no concurrency inside a node, while threads residing at a different node may run concurrently. The daemon also guarantees non-interruptible access to each link. While a thread is executing a certain link's method, no other threads can access the same link. The thread will suspend its non-interruptible execution and relinquish CPU to a next ready thread by calling one of the following methods

- Network construction/destruction affecting a remote daemon
  - *[e]create node* if it must be created on a remote daemon
  - *[e]destroy node*

    - *[e]create link* if it must be established to a remote daemon

- Link-accessing statements
    - *link*[*LinkID*]*.member*
    - *link*[*LinkID*]<*LinkClass*>*.member*
- All thread deployment statements
    - [e]hop
    - [e]hopalong
    - [e]fork
    - [e]forkalong
- Inter-threads synchronization
    - *sleep*
    - *wakeup/wakeupall* if you have chosen *pthread* in your profile.

Note that an M++ thread never yields CPU unless it calls one of the above statements. Programmers are responsible to prevent a thread from hogging CPU, (i.e., to prevent the other threads from starving for it.) If an M++ thread needs to wait for another thread, it should either execute *sleep* to suspends its execution or repeatedly invoke *hop( )* to relinquish CPU to another thread. (The former is recommended for better performance.)

In the following code example, the *WakeUp* thread creates five *PrintMsg* threads and performs a busy wait until they increment the *num* node variable, display a "Good night" message, and fall asleep. Thereafter, the *WakeUp* thread prints out "Wake up", and wakes up all those five threads that then reply "Good morning".

```
const int threadNum = 5;

class Place {                    // a class used to instantiate a node
public:
  Place( ) : num( 0 ) { }
  int num;
};

thread PrintMsg {                // a child thread created by a Wakeup thread
public:
  void main( ) {
    node<Place>.num++;

    cout << ``Good night.'' << endl;
    sleep node;
    cout << ‘‘Good morning.’’ << endl;
  }
}

thread WakeUp {                  // a parent thread to spawn 5 PringMsg threads
private:
  int i;
public:
  void main( );
}

void WakeUp::main( ) {
  create node<Place> with( 0 );
  hop( 0 );

  for ( i = 0; i < threadNum ; i++ )
    create thread<PrintMsg>;

  while ( node<Place>.num != threadNum )
    hop( );

  cout << ‘‘Wake up!!’’ << endl;
  wakeupall node;
}
```

The above program runs in the following sequence of thread context switches:

1. The *WakeUp* thread performs all thread instantiations without any intervention. Thereafter, it relinquishes CPU to one of the five *PrintMsg* threads by calling *hop( )*.
2. Once acquiring CPU, each *PrintMsg* thread increments *num*, prints out a "Good night" message, and thereafter executes *sleep* that yields CPU to another thread. Eventually, all the *PrintMsg* threads have fallen asleep, and CPU comes back to the *WakeUp* thread.
3. The *WakeUp* thread executes *wakeupall node* and terminates itself. Control goes to one of *PringMsg* threads.
4. The *PrintMsg* thread, upon displaying "Good night", terminates its execution and passes control to another thread that is about to print out the same message.

4.3.8 Node Clocking

This feature permits an M++ thread to invoke the same method of all nodes existing over the system. The method to be defined in node classes must have the following interface:

```
void clocking( int time );
```

To invoke all the *clocking* methods, an M++ thread must execute the following *clocking* statement in its *main( )* method.

- *clocking time;*
  passes an integer argument, *time* to the *clocking* method of all nodes.

With node clocking, a thread can avoid visiting all nodes over the system in order to invoke their *clocking* method, which may be useful for simulation applications to tick their virtual time.

Since a thread can access links while executing a method of the current node, (i.e., it can call link methods from a node method), the *clocking* method can be used to exchange data via links between any two neighboring nodes at a time.

Note that node clocking does not ideally mean a simultaneous invocation of all node's *clocking( )* method. In particular, *clocking( )* methods on the same daemon will be invoked in sequential. Of course, such calling sequence is not deterministic, either.

4.3.9 Communication

While threads share node variables on the same node and thus use them as indirect communication media, they can also establish direct communication by calling each other's public method. Note that a called thread must be sleeping on the same node where the calling node resides.

- *thread*[*ThreadID*]*.member*
  calls a predefined *member* function, (i.e., *name( )* or *id( )*) of the thread *sleeping* on the current node and specified in *ThreadID*.
- *thread*[*ThreadId*]<*ThreadClass*>*.member*
  calls a user-defined *member* function of the thread *sleeping* on the current node and specified in *ThreadID*.

The following program includes two threads such as *Container* and *Reader*. The former receives one argument upon its initialization and maintains it as an integer value. The latter, through the inter-thread communication, reads this integer value of each *Container* thread residing on the current node. More specifically, *Reader* scans all threads whose identifications can be found from *thread.min( )* to *thread.max( )*, checks if each of them has "Container" as its class name, and, if so, calls its *getData( )* method to retrieve the *number* value.

```
thread Container { // a thread keeping an integer value which will be read by Reader
public:
  Container( int argc, const char** argv ) : number( atoi( argv[4] ) ) { }
  int getData( ) { return number; }
  void main( ) { sleep; }
private:
  int number;
};

thread Reader {    // a thread reading each Container's integer value
public:
  void main( );
private:
  thread_id i;
}

void Reader::main( ) {
  if ( thread.count( ) > 0 ) {
    for ( i = thread.min( ) ; i <= thread.max( ) ; i = thread.next( i ) )
      if ( thread[i].name( ) == ``Container'' )
        cout << ``Container['' << i << ``]'s data = ''
             << thread[i]<Container>.getData( ) << endl;
  }
}
```

4.3.10 Global Virtual Time (Soon Made Available)

This feature has been supported at the daemon level but not yet at the language level. It will be soon made available.

Using the M++ global virtual time (abbreviated as GVT), threads can schedule their next action at a certain logical time that starts from time 0. This feature is intended to ease discrete-event simulations with M++. The following five methods are used to start, schedule thread actions along, and stop GVT.

- *void gvt_ready( )*
  A thread which wants to participate in a GVT computation must call *gvt_ready* in prior to calling any of the following methods: *[e]hop*, *[e]fork*, *[e]create*, *[e]destroy* and all GVT methods except *gvt_start*.

  When calling *gvt_ready* at time 0, the thread waits until another thread calls *gvt_start* that triggers time ticking. If *gvt_start* has been already called and a

GVT computation is in progress, the thread does not suspend its execution and goes forward to the next *gvt_delta* function call.

- *void gvt_start( )*
  triggers GVT ticking from 0 and wakes up all threads which have called *gvt_ready*. A thread calling *gvt_start* implicitly participates in a GVT computation and thus does not have to call *gvt_ready*.
- *void gvt_delta(int dTime)*
  Once participating in a GVT computation through *gvt_ready* or *gvt_start*, a thread may call *gvt_delta* to suspend its execution until GVT reaches the current time plus *dTime*.
- *int gvt_time( )*
  returns the current virtual time which is 0 or a positive integer.
- *void gvt_end( )*
  gets out of the current GVT computation. Each participating thread implicitly leaves the GVT computation upon its termination. However, if a thread needs to behave regardless of GVT, it must explicitly call *gvt_end*.

Note that GVT does not tick up as far as any participating threads are still active. Once all of them have called *gvt_delta* to suspend themselves, the system increments GVT to the nearest future time at which at least one thread schedules its next action.

5 Examples

This section deals with two entire sets of M++ code such as Section 5.1: Gas Station and Section 5.2: Ant Farm. The former is a uniprocessor example program in that *SportsCar* and *Truck* agents stop by a gas station to be pumped by a *GSCashier* agent. The latter is a multiprocessor application to find the gene of the most efficient ant forager by walking ants over a meshed farm constructed over multiple workstations.

5.1 Gas Station

This simple program models a gas station cashier fueling two different typed vehicles such as sports-cars and trucks. As shown in Figure 3, the former receives regular oil and the latter diesel.

Those three simulation entities are programmed as a *Cashier*, a *SportsCar*, and a *Truck* thread respectively. In particular, the latter two threads are derived from the

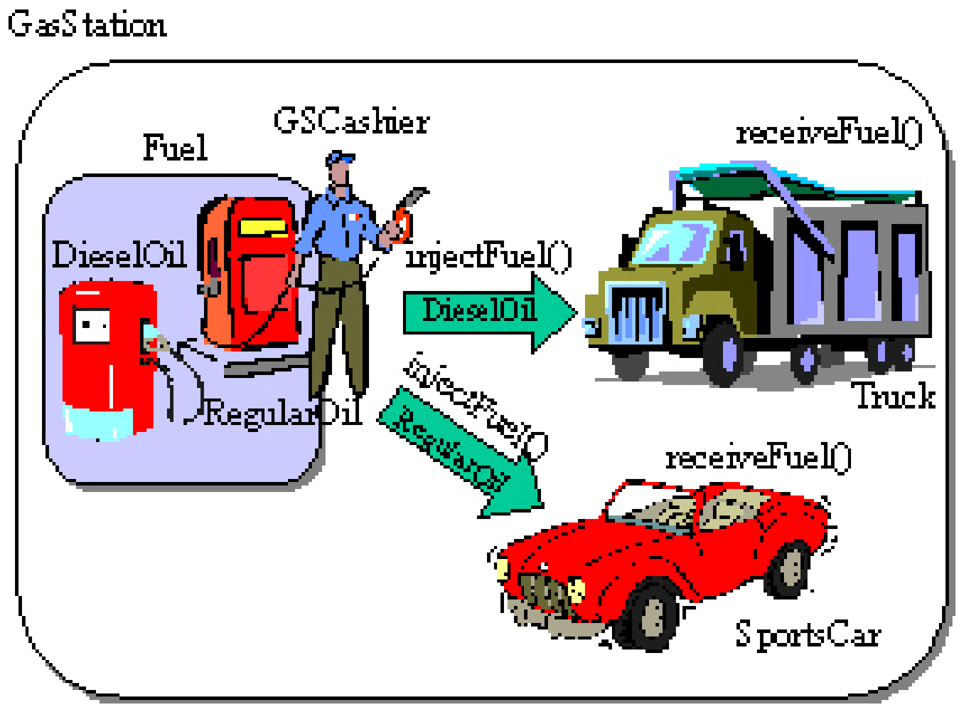

Figure 3: Cashier and Cars

Car base thread. In addition to those threads, the program defines the *GasStation* and the *Fuel* class, each instantiated as where *Cashier* works and what it serves.

5.1.1 GasStation/Fuel Classes

The *gasStationId* constant integer defines the id of the node where a *Cashier* thread works. It simply specifies 0.

The *GasStation* class includes nothing. It is used for a *Cashier* thread to instantiate its working node (whose id is *gasStationId*, i.e., 0).

The *Fuel* class includes two private integer variables: *dieselOid* and *regularOil*, each having 1000 (gallons). The class provides four public methods:

Methods	Descriptions
int useDieselOil(amount)	subtract amount or all from diesel and return this amount
int useRegularOil(amount)	subtract amount or all from regular and return this amount
void print( )	display the amount of diesel/regular oil left
bool empty( )	return true if both diesel and regular are empty

```
const int gasStationId( 0 );     // node id

class GasStation                 // node class
{
};

class Fuel                       // maintained by GSCashier
{
public:
  Fuel() : dieselOil( 1000 ), regularOil( 1000 ) {}
  int useDieselOil( int amount ) // subtract amount if possible, otherwise
  {                              // subtract all from diesel diesel oil
    if( dieselOil < amount )
      amount = dieselOil;

    dieselOil -= amount;
    return( amount );
  }

  int useRegularOil( int amount ) // subtract amount if possible, otherwise
  {                               // subtract all from regular oil
    if( regularOil < amount )
      amount = regularOil;

    regularOil -= amount;
    return( amount );
  }

  void print()                    // print out the amount of diesel/regular
    { cout << "Diesel: " << dieselOil << "  Regular: " << regularOil << endl; }

  bool empty()                     // true if no oil remains.
    { return( dieselOil == 0 && regularOil == 0 ); }

private:
  int dieselOil;
  int regularOil;
};
```

5.1.2 Car/SportsCar/Truck Threads

The *Car* thread is the base of both *SportsCar* and *Truck* threads. It includes two integers: *fuel* indicating the fuel remaining in its tank and *tankmax* specifying the tank capacity. Both are initialized upon a constructor invocation. This base thread has two public methods:

Methods	Descriptions
void injectFuel(Fuel& fuel)	pure virtual function implemented by *SportsCar* and *Truck*
void receiveFuel( )	wake up and wait for a *GSCashier* thread that pumps oil and signals it back

Note that, since *Car* is an abstract thread, it does not have the *main( )* method.

```
thread Car
{
public:
  Car( int tankMax ) : tankMax( tankMax ), fuel( 0 ) {}
  virtual void injectFuel( Fuel& fuel ) = 0;  // implemented by each sub thread
protected:
  void receiveFuel()  // find and wait for GSCashier to complete pumping gas
    {
      if( thread.count() )
        for( i = thread.min() ; i <= thread.max() ; i = thread.next( i ) )
          if( thread[i].name() == "GSCashier" )
            {
              wakeup i;   // wake up GSCashier
              sleep;      // sleep till GSCashier pump fuel and wake me up
              break;
            }
    }

  int fuel;          // the current amount of fuel stored in the tank
  int tankMax;       // the tank capacity
  thread_id i;       // the id of GSCashier thread
};
```

The *Truck* thread implements the *injectFuel* method that actually receives a *Fuel* object from *GSCashier* and calls its *useDieselOil( )* method. Through *useDieselOil*, the *Truck* thread fills its diesel tank.

The *main( )* method, (i.e., *Truck*'s behavioral scenario), migrates *Truck* to where *GSCashier* is working; wakes up and waits for *GSCashier* by calling *receiveFuel( )* method; and resumes its execution after *GSCashier* calls *injectFuel( )* to pump fuel and notify it.

```
thread Truck : public Car
{
public:
  Truck( int argc, const char** argv ) : Car( 100 ) {}
  void injectFuel( Fuel& fuel )              // called by GSCashier
    { Car::fuel = fuel.useDieselOil( tankMax - Car::fuel ); }
  void main()
    {
      hop( gasStationId );
      cout << "Truck: I'm at the GasStation." << endl;
      receiveFuel();                         // wait till GSCashier pumps fuel
      if( fuel == tankMax )
        cout << "Truck: I'm filled with fuel." << endl;
    }
};
```

The *SportsCar* thread implements the *injectFuel* method that actually receives a *Fuel* object from *GSCashier* and calls its *useRegularOil( )* method. Through *useRegularOil*, the *SportsCar* thread fills its regular-oil tank.

Its *main( )* method is the same as that of *Truck* except some message-printing statements.

```
thread SportsCar : public Car
{
public:
  SportsCar( int argc, const char** argv ) : Car( 50 ) {}
  void injectFuel( Fuel& fuel )
    { Car::fuel = fuel.useRegularOil( tankMax - Car::fuel ); }
  void main()
    {
      hop( gasStationId );
      cout << "SportsCar: I'm at the GasStation." << endl;
      receiveFuel();
      if( fuel == tankMax )
        cout << "SportsCar: I'm filled with fuel." << endl;
    }
};
```

5.1.3 GSCashier

The *GSCashier* thread creates a new node with its *gasStationId* (= zero), hops to it, and repeats the rest of its behavior enclosed in its *while* loop. It sleeps until woken by either *Truck* or *SportsCar*. Thereafter, it identifies which vehicle thread has woken it up and calls this vehicle's *injectFuel* method, as passing its *Fuel* object, which pumps the appropriate type of oil to the vehicle. Finally, *GSCashier* wakes up the vehicle thread and displays its *fuel* information.

```
thread GSCashier
{
public:
  void main()
    {
      create node<GasStation> with( gasStationId ); // create a working place
      hop( gasStationId );                          // work there

      while( 1 )
        {
          cout << "GSCashier: Waiting for cars." << endl;
          sleep;                   // wait till a new vehicle comes around

          if( thread.count() )
            for( i = thread.min() ; i <= thread.max() ; i = thread.next( i ) )
              {
                thread[i]<Car>.injectFuel( fuel );
                cout << "GSCashier: Injected fuel to "
                     << thread[i].name() << endl;
                wakeup i;

                fuel.print();
                if( fuel.empty() )
                  cout << "Fuel is empty." << endl;
              }
        }
    }
private:
  thread_id i;
  Fuel fuel;
};
```

All the source code was programmed in */home/m++/appl/m++/Car/Car.mpp.* To compile and execute it, follow the instruction given in Section 3.2.

5.2 Ant Farm

As shown in Figure 4, this multiprocessor application finds the most efficient ant forager by distributing food on a meshed network and walking many ant threads, each carrying different behavioral gene. To perform such a simulation, the program includes the following five types of M++ threads:

1. *Mesh* constructs a given size of mesh network.

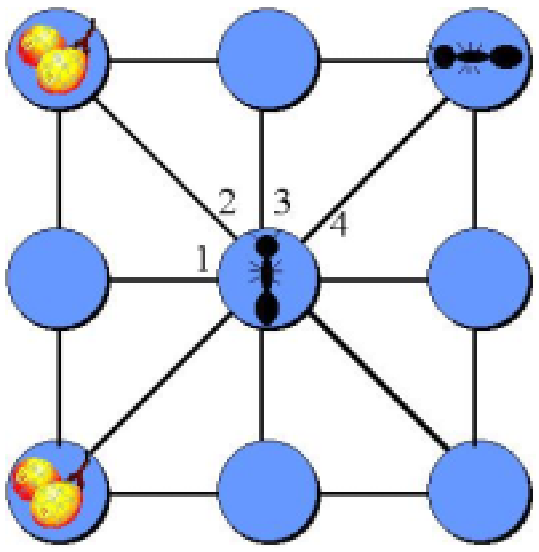

Figure 4: Ant farm

2. *InitEnvironment* distributes food on four cells, each located at diagonally one inside a different corner of the mesh network.
3. *Ant* searches for and bring back food to the nest.
4. *Gvt* performs inter-threads synchronization every tick of simulation time.
5. *CheckResult* shows food, pheromone and ant distribution over the network after simulation.

The following subsections list each of those thread programs.

5.2.1 DaemonVar/Node/Link Definitions

Those classes are used to construct a *farm* simulation space where *ant* agents walk around.

DaemonVar
This is a *daemonobj* class mainly used for two functionalities such as virtual time maintenance and node-to-daemon mapping. Note that we developed this ant farm application before completing the *M++* global virtual time maintenance feature that therefore had to be implemented at an application level. The following summarizes the variables defined in *DaemonVar*.

type	name	used for	descriptions
int	daemonTotal	node mapping	# of daemons over the system
int	meshsize	node mapping	a *meshsize*^{2} farm created
int	divider	node mapping	dividing farm in 1: strip or 0: lattice
int	ant_num	time maintenance	# of ants on this daemon
int	live_ant	time maintenance	# of active ants on this daemon
int	sleep_ant	time maintenance	# of sleeping ants on this daemon
int	gvt	time maintenance	the current global virtual time
bool	is_reported	time maintenance	false if none of daemons has informed the master of ant info.
bool	gvt_kicker_flag	time maintenance	false if this daemon has not informed the master of ant info.
Timer	timer	performance	see section **??**

```
import <cmath> {}
import <cstdlib> {}
import "Timer.h" { class Timer; }   // used for performance evaluation

const bool DEBUG( false );

class DaemonVar
{
public:
    DaemonVar() {
      ant_num = 0;
      live_ant = 0;
      sleep_ant = 0;
      gvt = 0;
      is_reported = false;
      gvt_kicker_flag = false;
      divider = 0;
    }
    int daemonTotal;      // # daemons over the system
    int meshsize;         // a meshsize x meshsize farm created
    int divider;          // dividing farm in 1: stripe or 0: lattice
    int ant_num;          // # ants residing on this daemon
    int live_ant;         // # active ants on this daemon
    int sleep_ant;        // # sleeping ants on this daemon
    int gvt;              // the current virtual time
    bool is_reported;     // false if none of daemons informed master of ants
    bool gvt_kicker_flag; // false if a daemon informed master of ants info.
    Timer timer;          // used for performance evaluation

    int getNodeID( int x ,int y ) { // decide a node id from [x,y] coordinates
      return y * meshsize + x;
    }
    int getDaemonID( int nodeId ) { // decide a daemon id from a given node id
      int dId;

      dId = 0;
      if ( divider == 1 )           // stripe mapping
```

```
  dId = ((nodeId * daemonTotal) /(meshsize*meshsize));
else {                        // lattice mapping
  switch( daemonTotal ) {
    case 1:
      dId = 0;
      break;
    case 2:
      if ( nodeId < meshsize * meshsize / 2 )
        dId = 0;
      else
        dId = 1;
      break;
    case 4:
      if ( nodeId < meshsize * meshsize / 2 ) {
        if ( nodeId % meshsize < meshsize / 2 )
          dId = 0;
        else
          dId = 1;
      } else {
        if ( nodeId % meshsize < meshsize / 2 )
          dId = 2;
        else
          dId = 3;
      }
      break;
    case 8:
      if ( nodeId < meshsize * meshsize / 4 ) {
        if ( nodeId % meshsize < meshsize / 2 )
          dId = 0;
        else
          dId = 1;
      }
      else if ( nodeId < meshsize * meshsize / 2 ) {
        if ( nodeId % meshsize < meshsize / 2 )
          dId = 2;
        else
          dId = 3;
      }
      else if ( nodeId < meshsize * meshsize / 4 * 3) {
        if ( nodeId % meshsize < meshsize / 2 )
          dId = 4;
        else
          dId = 5;
      } else {
        if ( nodeId % meshsize < meshsize / 2 )
          dId = 6;
        else
          dId = 7;
      }
```

```
            break;
          default:
            dId = 0;
            break;
          }
        }
        return dId;
     }
};
```

Trail
This is a *node* class that maintains geometric information and food. Each node is first created by a *Mesh* thread, is thereafter initialized by a *InitEnvironment*, is used for *ant* threads to communicate with one another, and is finally checked of its data by a *CheckResult* thread.

```
class Trail {
public:
    /*****************
    //Class Variable
    *****************/
    int x,y;              // coordinates
    int nest;             // 1 if it is the nest, otherwise 0
    int food;             // the amount of food left on this node
    int ant_num;          // # of ants on this node
    int pheromon_to_food; // the level of food-directive pheromone left on this node
    int pheromon_to_nest; // the level of nest-directive pheromone left on this node

    /*****************
    //Method
    *****************/

    Trail(int x,int y):x(x),y(y){   // constructor
        nest=0;
        food=0;
        ant_num=0;
        pheromon_to_food=0;
        pheromon_to_nest=0;
    }

    void drop_pheromon_to_food(int ph){   // increment food-directive pheromone
        pheromon_to_food = pheromon_to_food + ph;
    }

    void drop_pheromon_to_nest(int ph){    // increment nest-directive pheromone
        pheromon_to_nest = pheromon_to_nest + ph;
    }
    int get_pheromon_to_food(){            // detect food-directive pheromone
        return pheromon_to_food ;
    }

    int get_pheromon_to_nest(){            // detect nest-directive pheromone
        return pheromon_to_nest ;
    }

    void ant_in(){                         // one ant came in this node
        ant_num++;
    }

    void ant_out(){                        // one ant left this node
        ant_num--;
    }

    int is_food(){                         // return the amount of food
        return (food > 0);
    }

    int  get_food(int f){                  // return a requested or all
        if ( food < f )                    // remaining food
          f = food;
        food = food - f;
        return f;
    }

    void set_food(int f){                  // initialize the amount of food
        food=f;
    }

    void drop_food(int f){                 // receive food dropped by an ant
        food=food + f;
    }

    int is_nest(){                         // 1 if this is the nest
        return nest;
    }
};
```

Way

This is used to instantiate *link*s. The current implementation does not use any methods defined in *Way*.

```
class Way {
public:
    /*****************
    //Class Variable
    *****************/
    int i;
    int pheromon_to_food;
    int pheromon_to_nest;

    /*****************
    //Method
    *****************/
    Way(int i):i(i) {}

    /* Ant threads call none of those functions in the current implementation */
    void drop_pheromon_to_food(int ph){
        pheromon_to_food = pheromon_to_food + ph;
    }

    void drop_pheromon_to_nest(int ph){
        pheromon_to_nest = pheromon_to_nest + ph;
    }

    int get_pheromon_to_food(){
        return pheromon_to_food ;
    }

    int get_pheromon_to_nest(){
        return pheromon_to_nest ;
    }
};
```

5.2.2 Mesh Thread

The *Mesh* thread obtains a mesh size and a node mapping scheme from the first and the second argument respectively. For instance, given 200 as the mesh size and 1 as the node mapping scheme, it creates a 200 × 200 mesh network, divides it in stripe, and maps each divided space to a different processor. If a user wants to divide the network in lattice, the second argument must be set to 0.

The *Mesh* thread carries following six thread variables with it:

Variables	Descriptions
meshsize	a mesh size
divider	a node mapping scheme (0 = lattice, 1 = stripe)
i, j	coordinates of each *Trail* node
daemonNO	the daemon ID computed from a given node ID
ToNodeID	the node ID computed from given *i,j* coordinates

In addition, the *mesh* thread has the following two methods:

- *createMesh* first creates node[i,j] where $0 <= i < meshsize$ and $0 <= j < meshsize$, and then call its *link_create* method to establishes links emanating from each node.
- *link_create* creates a link from a given node to each of 8 neighboring nodes.

The *Mesh* thread is the first one injected into the system, which in turn means that *Mesh* is in charge of initializing all *daemonobj* objects. The following lists a complete code of the *Mesh* thread.

```
daemonobj DaemonVar();          // initialize daemonobj at each daemon

thread Mesh
{
public:
    Mesh( int argc, const char** argv ) :
       meshsize( atoi( argv[4] ) ), divider( atoi( argv[5] ) ) {}
    void main();                // argv[4]: meshsize, argv[5]: divider
private:
    int meshsize;               // a mesh size
    int i,j;                    // coordinates of each Trail node
    int daemonNO;               // the daemon ID computed from a given node ID
    int ToNodeID;               // the node ID computed from given i,j coordinates
    int divider;                // a node mapping scheme (0=lattice, 1=stripe)

    void link_create();
    void createMesh();
};

void Mesh::link_create()        // creates a link to each of 8 neighbors.
{
    // if not on y axis, create a link to west(0)
    if( node<Trail>.x != 0 ){
        ToNodeID = node.id() - meshsize;
        create link<Way>( 0 ) with( 0 )
          to( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) ) with( 4 );
    }
    // if not on y axis nor on the upper boundary, create a link to northwest(1)
    if( node<Trail>.x != 0 && node<Trail>.y != meshsize - 1 ){
        ToNodeID = node.id() - meshsize + 1;
        create link<Way>( 0 ) with( 1 )
          to( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) ) with( 5 );
    }
    // if on the upper boundary, create a link to north(2)
    if( node<Trail>.y != meshsize - 1 ){
        ToNodeID = node.id() + 1;
        create link<Way>( 0 ) with( 2 )
          to( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) ) with( 6 );
    }
    // if not on the upper nor the left boundary, create a link to northeast(3)
    if( node<Trail>.y != meshsize - 1 && node<Trail>.x != meshsize - 1 ){
        ToNodeID = node.id() + meshsize + 1;
        create link<Way>( 0 ) with( 3 )
          to( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) ) with( 7 );
    }
}

void Mesh::createMesh()         // create a [i,j] node and establish links from it.
{
    for(i=0;i<meshsize;i++){
        for(j=0;j<meshsize;j++){
            ToNodeID = meshsize * i+j;   // compute node id from i,j coordinates
            create node<Trail>( i, j )   // create a node
              with( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );
       }
    }

    for(i=0;i<meshsize;i++){
        for(j=0;j<meshsize;j++){
            ToNodeID = i*meshsize+j;
            hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID )); // hop to each node
            link_create();               // create links from it
        }
    }
}

void Mesh::main()
{
    for( i = 0 ; i < daemon.total() ; i++ ) {
      hop( INIT@i );                          // visit each daemon
      daemonobj.daemonTotal = daemon.total(); // initialize daemonobj
      daemonobj.meshsize = meshsize;
      daemonobj.divider = divider;
    }

    createMesh();                             // create a mesh network

    for ( i = 0 ; i < daemon.total() ; i++ ) {// revisit each daemon
      hop( INIT@i );                          // print out its task completion
      cout << "daemonobj.meshsize = " << daemonobj.meshsize << endl;
      cout << "daemonobj.divider = " << daemonobj.divider << endl;
    }
}
```

5.2.3 InitEnvironment Thread

The *InitEnvironment* thread receives one argument as the number of ants to inject, records it in each daemon's *daemonobj.ant_num*, and distributes food on four cells, each located at diagonally one inside a different corner of the mesh network. Given N as the mesh size, the cells to include food are: $cell[1,1]$, $cell[1,N-2]$, $cell[N-2,1]$, and $cell[N-2,N-2]$.

The *InitEnvironment* thread carries following three thread variables with it:

Variables	Descriptions
k	a for-loop control variable
ToNodeID	the node ID computed from given *i,j* coordinates
ant_num	a variable to store the total number of ants

The *InitEnvironment* thread initializes the simulation space using its three methods below:

- *setNest* sets the central cell of the network to the nest.
- *setFood* places 10,000 food on four cells, each located at diagonally one inside each corner of the mesh network
- *set_daemon_variable* have each daemon record the number of ants to be injected.

The following lists a complete code of the *InitEnvironment* thread:

```
thread InitEnvironment
{
public:
    InitEnvironment( int argc, const char** argv ) :
      ant_num( atoi( argv[4] ) ) {}       // argv[4]: #ants to be injected
    void main();
    void setNest();
    void setFood();
    void set_daemon_variable();
private:
    int k;                                // a for-iteration variable
    int ToNodeID;                         // a node id computed from i,j
    int ant_num;                          // #ants to be injected
};

void InitEnvironment::setNest()           // set the central cell to the next
{
    int halfsize = daemonobj.meshsize / 2;
    ToNodeID  = halfsize * daemonobj.meshsize + halfsize ;
    hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );
    node<Trail>.nest = 1;

    printf( "NEST LOCATION:(%d,%d)\n", node<Trail>.x, node<Trail>.y );
    printf( "IS NEST:%d\n", node<Trail>.is_nest() );
}

void InitEnvironment::setFood()           // place 1000 food each at four nodes
{
    //**** North East
    ToNodeID  =  daemonobj.meshsize + (daemonobj.meshsize-1) ;
    hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );
    node<Trail>.food=10000;

    //**** South East
    ToNodeID  = (daemonobj.meshsize-1)*daemonobj.meshsize
              + (daemonobj.meshsize-1);
    hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );
    node<Trail>.food=10000;

    //**** North West
    ToNodeID  = daemonobj.meshsize + 1;
    hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );
    node<Trail>.food=10000;

    //**** South West
    ToNodeID  = (daemonobj.meshsize-1 )*daemonobj.meshsize + 1;
    hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );
    node<Trail>.food=10000;
}

void InitEnvironment::set_daemon_variable()
{
    for ( k = 0; k < daemonobj.daemonTotal; ++k ){
      hop( INIT @ k );                 // visit each daemon
      daemonobj.ant_num = ant_num;     // # ants to be injected
      daemonobj.live_ant = ant_num;    // they are initially all active
      daemonobj.sleep_ant = 0;         // they are not sleeping of course
      daemonobj.gvt = 0;               // simulation starts from time 0
      daemonobj.is_reported = false;   // 1st synchronization has not been reported

      if( DEBUG )
        {
          printf("----------set_daemon_variable---------\n");
          printf("daemonobj.meshsize:%d\n",daemonobj.meshsize );
          printf("daemonobj.ant_num:%d\n",daemonobj.ant_num  );
          printf("daemonobj.live_ant :%d\n", daemonobj.live_ant);
          printf("daemonobj.sleep_ant:%d\n",daemonobj.sleep_ant );
          printf("daemonobj.gvt :%d\n", daemonobj.gvt);
        }
    }
}

void InitEnvironment::main()
{
  cout << "daemonobj.meshsize = " << daemonobj.meshsize << endl;
  cout << "daemonobj.divider = " << daemonobj.divider << endl;

  setNest();
  setFood();
  set_daemon_variable();

  cout <<"init environment END " <<endl;
}
```

5.2.4 Ant Thread

The *Ant* thread repeats the following actions every simulation cycle: (1) pick up food if it finds any, (2) drop off food in the nest or on route to the nest accidentally, (3) follow the strongest pheromone, (4) decide the next direction if it feels no pheromone, and (5) secrete some pheromone and walk to the next node. To give more reality to each action, the gene of each ant characterizes its average number of food to pick up, average pheromone amount to emit, pheromone sensitivity, and durability to walk straight. Those characteristic parameters are maintained in the following variables:

Variables	Descriptions
ToNodeID	the node ID computed from given *i,j* coordinates
food	the amount of food which this ant is currently holding
total_food	the total amount of food which this ant has carried to the nest
step	the current simulation time
direction	the next moving direction
ph_react	the current threshold of pheromone amount for ant to react.
i, j	for-loop control variables
border	a flag to indicate if the ant is on a processor border
randVar	the latest random number obtained from *rand_r( )*
gene	gene information obtained using *random( )*
norm_u	the average of this ant's pheromone reaction level
ph_react0	the initial pheromone reaction threshold
pois_l	the average amount of food this ant picks up
keep_p	the probability to keep holding the food which the ant picked up
straight_p	the probability to go straight
trust_p	the probability to trust the latest pheromone information
seed	this ant's seed used for *rand_r( )*
temp_node_id	the node id to make the ant return to this node later
temp_daemon_id	the daemon id to make the ant return to this daemon later
activate_pheromon	the flag to make the ant follow pheromone
temp_max	the variable to store pheromone considered maximal temporarily
temp_direction	the variable to store the next moving direction temporarily

The following methods describe each action for an *Ant* to take:

- *init* calls *set_gene* to initialize thread variables based on a given gene and migrates a thread to the nest node.
- *set_gene* randomly decides the genetic code of each ant and initializes thread variables such as *norm_u*, *ph_react0*, *pois_l*, *keep_p*, *straight_p*, *trust_p*, and *seed*.
- *synchronize* migrates this ant to the *INIT* node on daemon #0 if it is the first one to do so, injects a new *gvt* thread if it has not yet been injected, comes back to the original node, and sleeps for the *gvt* thread to wake up this ant.
- *init_at_trail* increments the number of ants residing on this node and adds 10 to its own *ph_react* value, (which makes this ant less sensitive to pheromone.
- *work_at_trail* gets food if the current node is not the nest, drops food if it is the nest, and accidentally drops off food on route to the nest.
- *feel_pheromon* replaces this ant's variables *tempMax* and *tempDir* with the maximal pheromone and its direction which this ant has temporarily detected.
- *decide_direction* investigates the maximal pheromone and its direction by walking through all the cells neighbors of the current cell and decides the next moving direction based on this investigation, if this ant trusts this pheromone information. If not, the ant randomly chooses the next direction to walk.

- *go_next_trail* drops its pheromone, decrements the number of ants residing on the current node, and actually migrates the ant to the next cell.
- *fact* computes the factorial of a given argument.
- *normal* computes the normal distribution, given a average parameter and a random value.
- *poisson* computes the poisson distribution, given a lumda parameter and a random value.

The following shows a complete code of the *Ant* thread:

```
thread Gvt;                               // Ant spawns a Gvt thread
const int LIFE_STEP( 4001 );              // simulation ends in time 4001
extern "C" int rand_r(unsigned int *seed); // rand_r is used.

thread Ant
{
public:
    Ant( int argc, const char** argv ) :
      activate_pheromon( ( atoi( argv[4] ) == 1 ) ? true : false ) {}
    void main();                          // if 1st arg is 1, pheromone is used.
private:
    int ToNodeID;    // the node ID computed from given i,j coordinates
    int food;        // the amount of food which this ant is holding
    int total_food;  // the total amount of food which this ant carried to the nest
    int step;        // the current simulation time
    int direction;   // the next moving direction
    int ph_react;    // the current threshold of pheromone amount for ant to react
    int i, j;        // for-loop control variables
    bool border;     // a flag to indicate if the ant is on a processor border

    //GENE
    short  gene;            // gene information obtained using random( )
    double norm_u;          // the ave. of this ant's pheromone reaction level
    int ph_react0;          // the initial pheromone reaction threshold
    double pois_l;          // the ave. amount of food which this ant picks up
    double keep_p;          // the probability to keep holding food
    double straight_p;      // the probability to go straight
    double trust_p;         // the probability to trust the latest pheromone info.
    unsigned int seed;      // this ant's seed used for rand_r( )

    //To Report host0
    int temp_node_id;       // to make the ant return to this node later
    int temp_daemon_id;     // to make the ant return to this daemon later

    //IN decide direction
    bool activate_pheromon;  // the flag to make the ant follow pheromone
    int temp_max;            // to store pheromone considered maximal temporarily
    int temp_direction;      // to store the next moving direction temporarily

    //Poisson Normal
    double randVar;          // the latest random number obtained from rand_r()

    void init();
    void set_gene();
    void synchronize();
    void init_at_trail();
    void work_at_trail();
    void feel_pheromon( int& tempMax, int& tempDir, int& dir );
    void decide_direction();
    void go_next_trail();

    //Utility
    int fact( int i );
    int normal( double micro, double randVar );
    int poisson( double lumda, double randVar );
};

// Utility Function
int Ant::fact( int i )                           // computes the factorial of i
{
  int f;

  for ( f = 1; i > 0; i-- )
    f *= i;
  return( f );
}

int Ant::normal( double micro, double randVar )  // compute the normal distribution
{                                                // based on micro and randVar
  const float e = 2.718281828459;
  const float pi = 3.141592;
  double n = 0.0;
  int i;

  for ( i = 0;i<=10 ; i++ ) {
    n +=  1.0/sqrt( 2.0*pi ) * pow( e, -( pow( i - micro, 2.0 )/2.0 ) );
    if ( n >= randVar )
      break;
  }
  //cout << "normal=" << i << endl;
  return i;
}

int Ant::poisson( double lumda, double randVar ) // compute the poisson distribution
{                                                // based on lumda and randVar
```

```
  const double e = 2.718281828459;
  double p = 0.0;
  int i;

  for ( i = 0;i<=10 ; i++ ) {
    p += ( pow( e, -lumda ) * ( pow( lumda, i ) / fact( i ) ) );
    if ( p >= randVar )
      break;
  }
  //cout << "poisson=" << i << endl;
  return i;
}

void  Ant::set_gene( ){    // decide the ant gene

  int temp_gene;

  gene = random() & 0x0000ffff;
  temp_gene = gene & (0x1);

  switch (temp_gene) {     // decide the ave. of pheromone reaction level
  case 0:
    norm_u=3.0;
    break;
  case 1:
    norm_u=4.0;
    break;
  }

  temp_gene=gene >> 1;
  temp_gene=temp_gene & (0x3);

  switch (temp_gene) {     // decide the initial pheromone reaction threshold
  case 0:
    ph_react0=3;
    break;
  case 1:
    ph_react0=4;
    break;
  case 2:
    ph_react0=5;
    break;
  case 3:
    ph_react0=6;
    break;
  }

  temp_gene=gene >> 3;
  temp_gene=temp_gene & (0x1);
```

```
switch (temp_gene) {     // decide the ave. amount of food which ant picks up
case 0:
  pois_l=1.0;
  break;
case 1:
  pois_l=2.0;
  break;
}

temp_gene=gene >> 4;
temp_gene=temp_gene & (0x3);

switch (temp_gene) {     // decide the probability to keep holding food
case 0:
  keep_p=0.8;
  break;
case 1:
  keep_p=0.85;
  break;
case 2:
  keep_p=0.9;
  break;
case 3:
  keep_p=0.95;
  break;
}
temp_gene=gene >> 6;
temp_gene=temp_gene & (0x3);

switch (temp_gene) {     // decide the probability to go straight
case 0:
  straight_p=0.8;
  break;
case 1:
  straight_p=0.85;
  break;
case 2:
  straight_p=0.9;
  break;
case 3:
  straight_p=0.95;
  break;
}
temp_gene=gene >> 8;
temp_gene=temp_gene & (0x3);

switch (temp_gene) {     // decide the probability to trust the latest pheromone
case 0:
    trust_p=0.8;
    break;
  case 1:
    trust_p=0.85;
    break;
  case 2:
    trust_p=0.9;
    break;
  case 3:
    trust_p=0.95;
    break;
    }

  seed = ( unsigned int )gene;  // used for rand_r( )

  if( DEBUG )
    {
      printf("norm_u:%f\n",norm_u);
      printf("ph_react0:%d\n",ph_react0);
      printf("pois_l:%f\n",pois_l);
      printf("keep:%f\n",keep_p);
      printf("straight:%f\n",straight_p);
    }
}

void Ant::init()           // initialize the thread variables
{
  food=0;                  // no food held
  step=0;                  // time is 0
  direction=-1;            // no direction decided
  total_food=0;            // no food carried
  ph_react=0;              // the initial pheromone reaction level is 0
//Set Gene
  set_gene( );             // set my genetic information

//move to nest
  ToNodeID  = (daemonobj.meshsize/2)*daemonobj.meshsize
            + (daemonobj.meshsize/2) ;
  hop( ToNodeID @ daemonobj.getDaemonID( ToNodeID ) );  // hop to the nest
}

void Ant::synchronize()                     // visit daemon 0 to inject gvt
{
  if ( daemonobj.sleep_ant  ==
       ( ( daemonobj.ant_num > daemon.total( ) ) ?
 ( daemonobj.ant_num / daemon.total( ) - 1 ) : 0 )
       && daemonobj.gvt_kicker_flag == false ) {

    // I AM  GVT  Kicker
    daemonobj.gvt_kicker_flag=true;

    //Address to return to
    temp_node_id= node.id();
    temp_daemon_id=daemon.id() ;

    hop( INIT @ 0 );

    if( daemonobj.is_reported == false ) { // No one has injected gvt yet.
      daemonobj.is_reported = true;
      create thread<Gvt>;
    }
    hop( temp_node_id @ temp_daemon_id );  // go back where I was
  }

  daemonobj.sleep_ant++;                   // sleep there until gvt wakes me up
  sleep daemon;
}

void Ant::init_at_trail()                  // the first thing at a new node
{
  node<Trail>.ant_in( );
  ph_react += ( daemonobj.ant_num/10 );
}

void Ant::work_at_trail()                  // the main work at a new node
{
  // get food
  if( node<Trail>.is_food()
      && (node<Trail>.is_nest() != 1 ) ){

    randVar = fmod( rand_r( &seed ), 10000.0 ) / 10000.0;
    food = node<Trail>.get_food( poisson(pois_l, randVar) );
  }

  // drop food at nest
  if( node<Trail>.is_nest() == 1 ){
    if( food > 0 ) {
      node<Trail>.drop_food(food);
      total_food += food;
      food = 0;
    }
  }

  // drop food on route to nest
  randVar = fmod( rand_r( &seed ), 10000.0 ) / 10000.0;
  if( ( randVar > keep_p ) && food >= 1 ){
    node<Trail>.drop_food( 1 );
    food = food - 1;
  }
}

void Ant::feel_pheromon( int& tempMax, int& tempDir, int& dir )
{                                         // detect maximal pheromone and its
  int link_ph;                            // direction

  if ( food == 0 )
    link_ph = node<Trail>.get_pheromon_to_nest( );
  else
    link_ph = node<Trail>.get_pheromon_to_food( );

  if ( link_ph > tempMax ) {
    tempMax = link_ph;
    tempDir = dir;
  }
}

void Ant::decide_direction()     // decide which of eight links to hop along
{
    if( step == 0 )              // at time 0
      direction=rand_r( &seed ) % 8;   // decide the direction randomly

    temp_max=0;
    temp_direction=0;

    if ( activate_pheromon ) {   // must follow pheromone
      /* checking if neighbors exist remotely or do not exist */
      border = false;
      for (i = 0; i <= 7; i += 2 ) {
        if ( !( link.exists(i) ) ) {   // check if link i does not exist
          if ( i == 0 && node<Trail>.y == 0 ) {
            //cout << "node[0,0]!!!" << endl;
            continue;
          }
          //cout << "link=" << i << " not exist" << endl;
          border = true;
          break;
        }
      }
      if ( i > 7 ) {                   // check if link i goes to a remote node
        for ( i = 0; i <=7; i += 2 ) {
          if ( link[i].dest_daemon_id() != daemon.id( ) ) {
            break;
          }
        }
      }
      i = ( i + 7 ) % 8;
      /* searching  pheromon */
      if ( !border ) {                 // I am not on the border
        j=i;                           // detect each neighbor's pheromone
        hopalong( i );
        feel_pheromon( temp_max, temp_direction, j );
        ++j %= 8 ;
        hopalong( i = ( i + 3 ) % 8 );
        feel_pheromon( temp_max, temp_direction, j );
        ++j %= 8;
        hopalong( i );
        feel_pheromon( temp_max, temp_direction, j );
        ++j %= 8;
        hopalong( i = ( i + 2 ) % 8 );
        feel_pheromon( temp_max, temp_direction, j );
        ++j %= 8;
        hopalong( i );
        feel_pheromon( temp_max, temp_direction, j );
      } else {                         // I am on the border
i = ( i + 3 ) % 8;             // investigate only valid neighbors'
if ( link.exists(i) ) {        // pheromone
  border = false;
  j=i;
  hopalong( i );
          feel_pheromon( temp_max, temp_direction, j );
  ++j %= 8;
  hopalong( i = ( i + 2 ) % 8 );
          feel_pheromon( temp_max, temp_direction, j );
} else {
  border = true;
          i = ( i + 2 ) % 8;
          j=i;
          hopalong( i );
          feel_pheromon( temp_max, temp_direction, j );
        }
      }
      ++j %= 8;
      hopalong( i = ( i + 2 ) % 8 );
      feel_pheromon( temp_max, temp_direction, j );

      if ( !border ) {
        ++j %= 8;
        hopalong( i );
        feel_pheromon( temp_max, temp_direction, j );
      }
      ++j %= 8;
      hopalong( i = ( i + 2 ) % 8 );
      feel_pheromon( temp_max, temp_direction, j );

      ++j %= 8;
      hopalong( ( i + 2 ) % 8 );
    }

    randVar = fmod( rand_r( &seed ), 10000.0 ) / 10000.0;   // get a random number
    if ( randVar > trust_p || temp_max < ph_react ) {
      /* I don't trust or can't feel pheromone */
      do {
        /* I doubt my direction */
        direction = rand_r( &seed ) % 8;
        randVar = fmod( rand_r( &seed ), 10000.0 ) / 10000.0;
      } while ( randVar > straight_p );
    } else {
      /* I'll follow the pheromone */
      direction = temp_direction;
    }

    /* I'm on edge */
    if(node<Trail>.x==0)
      direction=4;
    else if(node<Trail>.x==daemonobj.meshsize-1)
      direction=0;
    else if(node<Trail>.y==0)
      direction=2;
    else if(node<Trail>.y==daemonobj.meshsize-1)
      direction=6;
}

void Ant::go_next_trail()
{
  // outward
  if( food == 0 ){
    randVar=fmod(rand_r( &seed ),10000.0)/10000.0; // get a new random number
    node<Trail>.drop_pheromon_to_food( normal(norm_u, randVar) );
    node<Trail>.ant_out();      // unregister itself at the current node
    hopalong(direction);        // move to a next new node
  }

  // homeward
  if( food > 0 ){
    randVar=fmod(rand_r( &seed ),10000.0)/10000.0; // get a new random number
    node<Trail>.drop_pheromon_to_nest(normal(norm_u,randVar));
    node<Trail>.ant_out();      // unregister itself at the current node
    hopalong(direction);        // move to a next new node
  }
}

void Ant::main()                // the main ant behavior
{
  init();                       // initialize its thread variables
  //STEP
  for ( step = 0; step < LIFE_STEP; step++ ) {
    synchronize();              // synchronized with gvt
    init_at_trail();            // register itself at a new node
    work_at_trail();            // pick up and drop off food
    decide_direction();         // search for food
    go_next_trail();            // go to a next new node
  }

  //To stop Timer
  synchronize();
  node<Trail>.ant_in();

  if( DEBUG ) {
    cout << "ant[" << selfMigThr.getId( ) << "] got "
         << total_food << " amount of food"
         << " (norm_u=" << norm_u
         << " ph_react0=" << ph_react0
         << " pois_l=" << pois_l
         << " keep_p=" << keep_p
         << " straight_p=" << straight_p
         << " trust_p=" << trust_p
         << " seed=" << seed << ")"
         << endl;
  }
}
```

5.2.5 Gvt Thread

The *Gvt* thread is injected by one of *Ant* threads every tick of simulation time and performs inter-threads synchronization. It repeats visiting every *INIT* node, (i.e., every daemon) until all *Ant* threads fall asleep, and wakes them up when confirming that they have slept. This thread carries the following three thread variables with it to achieve such synchronization:

Variables	Descriptions
ant_sum	the total number of ants over the system
sum	the variable used to sum up the number of ants
i	a for-loop control variable

The *gvt* thread consists of the following two sub-method:

- *check_around* repeats visiting every *INIT* node and summing up the number of sleeping ants (=*sum*) until the number reaches the total number of ants

(=*ant_sum*).

- *wake_around* visits every *INIT* node to wake up all ants sleeping on that daemon.

The following list is a complete code of the *check_result* thread.

```
thread Gvt
{
public:
    Gvt( int argc, const char** argv ) : sum( 0 ) {}
    void main();
private:
    int ant_sum;                      // # of ants over the system
    int sum;                          // sum up # ants on each daemon
    int i;                            // a for-loop control variable

    void check_around();              // check if all ants are sleeping
    void wake_around();               // wake up all ants sleeping on each daemon
};

void Gvt::check_around()              // check if all ants are sleeping
{
  ant_sum = daemonobj.ant_num;        // get the total number of ants
  while( 1 ) {
    sum = 0;
    for ( i = 0; i < daemonobj.daemonTotal; ++i ) {
      hop( INIT @ i );                // visit each daemon
      sum += daemonobj.sleep_ant;     // sum up the number of ants sleeping there
    }
    if( sum == ant_sum )              // if all ants are sleeping, return to main()
      break;
  }
}

void Gvt::wake_around()               // wake up all ants sleeping on each daemon
{
    for ( i = 0; i < daemonobj.daemonTotal; ++i ) {
        hop( INIT @ i );              // visit each daemon
        if( daemonobj.gvt == 0 )      // if virtual time is 0 start the timer
            daemonobj.timer.start( );

        daemonobj.gvt++;              // increment the virtual time
        daemonobj.live_ant = daemonobj.sleep_ant; // all sleeping ants will be woken up.
        if ( daemonobj.gvt % 500 == 0 ) { // show #ants at this daemon even 500 cycles
          cout << "gvt=" << daemonobj.gvt
               << ", ants=" << daemonobj.live_ant << endl;
        }
        daemonobj.sleep_ant = 0;      // no more sleeping ants
        daemonobj.is_reported = false;// synchronization variables are reset
        daemonobj.gvt_kicker_flag = false;
        wakeupall daemon;             // finally wake up all ants!
    }
}

void Gvt::main()
{
  if( daemonobj.gvt == LIFE_STEP ) {  // LIFE_STEP = 4001 in this implementation
      daemonobj.timer.stop();         // stop the timer
      cout <<"performance" <<endl;    // print out the performance data
      cout << daemonobj.timer.getElapsed( ) << endl;
      cout << daemonobj.timer.getInterval( ) << endl;
  }
  check_around();                     // otherwise, check if all ants are sleeping
  wake_around();                      // then, wake them up

  if ( daemonobj.gvt % 500 == 0 )     // print out the virtual time every 500 cycles
    cout << daemonobj.gvt <<endl;
}
```

5.2.6 CheckResult Thread

The *CheckResult* thread visits each node and displays its node status. This thread carries the following three thread variables with it:

Variables	Descriptions
i, j	the farm coordinates of each *Trail* node
ToNodeID	the node ID computed from given *i,j* coordinates

The *CheckResult* thread computes the node ID from each pair of *i,j* coordinates, hops to it, and displays its node variables: *x*, *y*, *food*, *pheromon_to_food*, *pheromon_to_nest*, and *ant_num*.

A complete code of the *check_result* thread is shown below:

```
thread CheckResult
{
public:
    void main();
private:
    int i,j;
    int ToNodeID;

    void checkMesh();
};
```

```
void CheckResult::checkMesh()
{
    for ( i = 0; i < daemonobj.meshsize; i++ ) {
      for ( j = 0; j < daemonobj.meshsize; j++ ) {
        ToNodeID = daemonobj.meshsize * i + j;
        hop(ToNodeID @ daemonobj.getDaemonID(ToNodeID));
        printf("mesh[%d,%d]: food=%d pheromon(ToNest=%d ToFood=%d) ants=%d\n",
               node<Trail>.x,
               node<Trail>.y,
               node<Trail>.food,
               node<Trail>.pheromon_to_food,
               node<Trail>.pheromon_to_nest,
               node<Trail>.ant_num
              );
      }
    }
}

void CheckResult::main()
{
  cout << "--------------------" << endl;
  checkMesh();
}
```

5.2.7 Compilation and Execution

The ant farm shown above was programmed in */home/m++/appl/m++/ant.mpp*. To compile this program, type:

```
% m++ ant.mpp
```

To run this program, run inject commands as follows:

```
% inject hostname ./Mesh 1 meshsize divider
% inject hostname ./InitEnvironment 1 #ants
% inject hostname ./Ant #ants 1
% inject hostname ./CheckResult 1
```

where *hostname*, *meshsize*, *divider*, and *#ants* have the following meaning:

parameters	remarks
hostname	the IP name of a computing node where each thread is injected
meshsize	the size of ant farm (meshsize * meshsize)
divider	1: divide ant farm in strip, or 0: divide ant farm in lattice
#ants	the total number of ants to be injected

Every time you start a new thread injection, we recommend you should wait for messages indicating that the previous injection and the subsequent thread action has completed.

6 Trouble Shooting

There are five possible errors you may encounter: abnormal termination, fopen error, bind system call error, segmentation fault/bus error, and dynamic module loading error. The following explains each error and its trouble shooting.

6.1 Abnormal Termination

An unexpected trap/interrupt or a Unix *kill* command may cause an abnormal termination and keep the previous socket connection in use. When it happens accidentally, you must wait for a few minutes until the underlying operating system recognizes that the previous socket port is already freed, or should edit your *profile* to change the port number.

6.2 Fopen error

This error occurs when an M++ daemon cannot locate *profile*. This file must be located in your current working directory where the daemon is invoked.

6.3 Bind system call error

This system call error occurs when a specified port number has been already allocated to socket communication. Possible scenarios and their respective solutions are:

1. You have already started an M++ daemon on this machine. Terminate the currently running daemon and restart a new daemon.
2. Someone is using the same port number to start his/her daemon. Redefine another port number in your *profile*.
3. Your previous M++ execution has resulted in an abnormal termination. Follow the instruction specified in Section6.1, "Abnormal Termination".

6.4 Segmentation fault/Bus error

If the *inject* command causes this error, the number of arguments actually given is incompatible to that of arguments coded in your thread program. You should issue this command with the correct number of arguments. If a segmentation fault occurs during computation, your M++ threads but not daemons may have caused this fault. You should carefully look through the header and source files of your thread programs.

6.5 Dynamic module loading error

If the system displays a message, *"SharedLib: Cannot load the dynamic module"*, an M++ daemon has resulted in failure of dynamic linking. There are two possible reasons to induce this error:

1. You might have given a wrong file to the *inject* command. In this case, simply re-inject your threads with an executable file name.
2. The daemon was unable to find out a class specified in *[e]create node* and/or *[e]create link* statements. Assure that you have defined a class corresponding to *ClassName* in the following statements:
   - *[e]create node<ClassName>(args_list) with (NodeID[@DaemonID]);*
   - *[e]create link<ClassName>(args_list) with (SrcID) to (NodeID[@DaemonID]) with (DestID);*

7 Final Comments

Focusing on multi-agent applications, the M++ system provides application designer with several important features: agent-based programming, runtime construction of simulation space, and virtual-time management. They can thus concentrate on modeling their multi-agent applications.

The M++ system is being tested by several users in University of Washington - Bothell, University of Tsukuba, University of California - Irvine, and Ehime University. M++ applications currently available include: ant farm, sugar mountain, remote exploration, and codi 1bit. Through more applications development and testing, we are planning to provide a debug environment, to enhance the system robustness, and implement dynamic load balancing.

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